

# **Proxy execution**

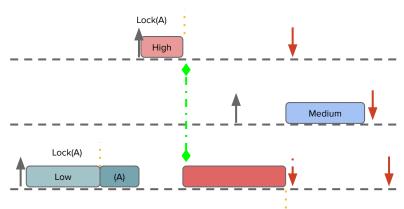
Juri Lelli <juri.lelli@redhat.com>

Linux Plumbers 2019



#### In a nutshell

- Priority inheritance mechanism
  - Applies to mutexes
  - Replaces (broken) deadline inheritance (SCHED\_DEADLINE)
    - Boosted task run outside runtime enforcement (!root prohibited)
  - Works across classes
  - Can unify mutex and rtmutex code
- Mutex owner can run using the scheduling context ("properties") of (several) donor(s)





### migrate\_task:

```
/*
* Follow blocked_on chain.
*/
for (p = next; p->blocked_on; p = owner) {
      mutex = p->blocked on;
      owner = __mutex_owner(mutex)
      if (task_cpu(owner) != task_cpu(p))
             * The blocked-on relation must not cross CPUs, if this happens
             * migrate @p to the @owner's CPU.
             *
             * This is because we must respect the CPU affinity of execution
             * contexts (@owner) but we can ignore affinity for scheduling
             * contexts (@p). So we have to move scheduling contexts towards
             * potential execution contexts.
```



## migrate\_task:

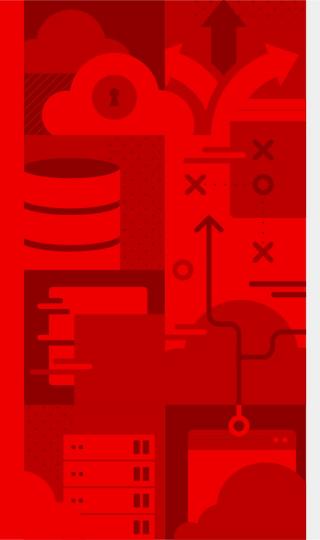
```
/*
* Follow blocked_on chain.
*/
                                                              Can we actually do this w/o
for (p = next; p->blocked_on; p = owner) {
                                                              breaking admission control?
      mutex = p->blocked on;
                                                              OSPM19 answer... NO!:-(
      owner = __mutex_owner(mutex)
      if (task_cpu(owner) != task_cpu(p))
             * The blocked-on relation must not cross CPUs, if this happens
             * migrate @p to the @owner's CPU.
             *
             * This is because we must respect the CPU affinity of execution
             * contexts (@owner) but we can ignore affinity for scheduling
             * contexts (@p). So we have to move scheduling contexts towards
             * potential execution contexts.
```



#### alternatives

- If p and owner allowed masks are equal -> migrate owner to p's cpu
  - If they were both admitted it means that their bw can be scheduled inside their root domain (DEADLINE doesn't currently care where)
  - What if owner is running? Wait until it is preempted?
  - But then maybe it's actually easier to let p (scheduling ctx) migrate to owner's cpu
    (as currently implemented)
- If they are disjointed neither p nor owner can be migrated to/from each other domain
  - Don't trigger proxy?
  - Is there theory already for tasks running on separate domains that share data?





# Thank you!

Juri Lelli <juri.lelli@redhat.com>

- in linkedin.com/company/red-hat
- youtube.com/user/RedHatVideos
- f facebook.com/redhatinc
- twitter.com/RedHat

