Moving swapping infrastructure to Rust

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About me: Vitaly

Konsulko Group

- Has been working with embedded Linux since 2003
 - MontaVista Europe
- Moved to Sweden in 2009
 - continued working with embedded
 Linux
 - consulting for Sony, Google, Axis, etc.
- Currently living in Malmö
- Working within the Konsulko Group
- Principal engineer at Konsulko AB



Why rustify?



- Bigger memory safety
 - Rust is not a silver bullet but still it has an edge over C there
 - NonNull, protection agains buffer overflows, etc.
- More secure
 - Reduced undefined behavior => smaller attack surface
- More robust
 - Emphasis on immutability
- More maintainable/sustainable
 - Powerful error handling mechanisms (e.g. Result, Option)
 - Strict coding style requiring comments



Kernel swapping structure



- Generic swap structure
 - Core swap implementation (swap.c, swapfile.c, etc.)
- RAM based compressed swap storages
 - Frontend (zswap, ZRAM)
 - Allocation backend (zsmalloc)
- Shmem (to a certain extent)
 - GPU drivers use it for GPU page swapping

Kernel swapping structure: what can be rustified?

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- Not all architectures support Rust for Linux [yet]
 - Core components should remain C based
 - zswap should remain C based
 - Shmem should remain C based
- zsmalloc could be a good candidate
 - More considerations will follow
- ZRAM is not a core component
 - Block device driver, basically a ramdisk on steroids
 - Used almost exclusively by x86_64 and aarch64

zsmalloc



- Seasoned allocator backend
 - Provides very good compression density
- Rather complex implementation
 - Objects are divided into 255 size classes
 - objects of the same class stored consequently within a page
 - Some objects span across 2 pages
- Doesn't work well in 16+K page setups
 - Lower granularity
- Currently the only available allocator backend
 - Can't be moved to Rust due to the architecture limitations

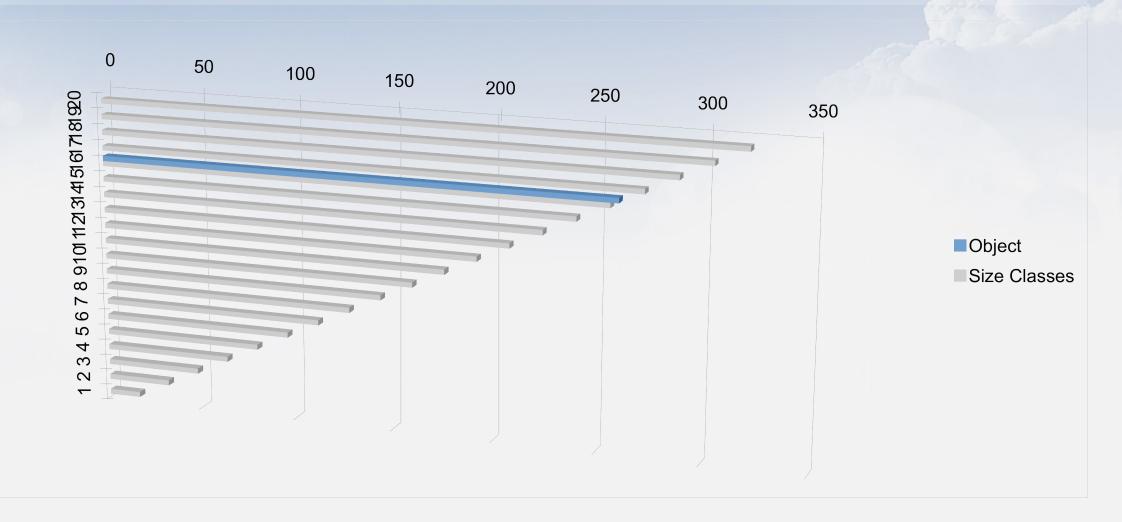
zblock



- If we can't move zsmalloc to Rust, we can implement an alternative allocator from scratch
- zblock: based on 2 simple ideas
 - Use page blocks instead of single pages
 - Divide large blocks into an array of same size objects
 - These same size objects (slots) don't have to be of 2^x size
- Small code footprint and easy to understand concept
- How it relates to zsmalloc
 - 4K pages: comparable compression density, better performance
 - 16K pages: better compression density, better performance

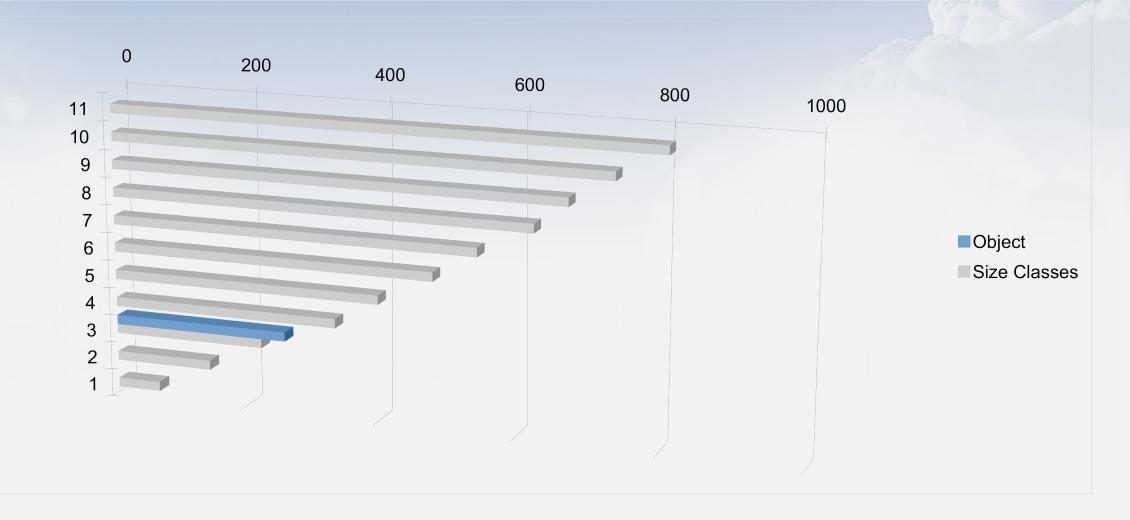
zsmalloc size classes





zblock size classes

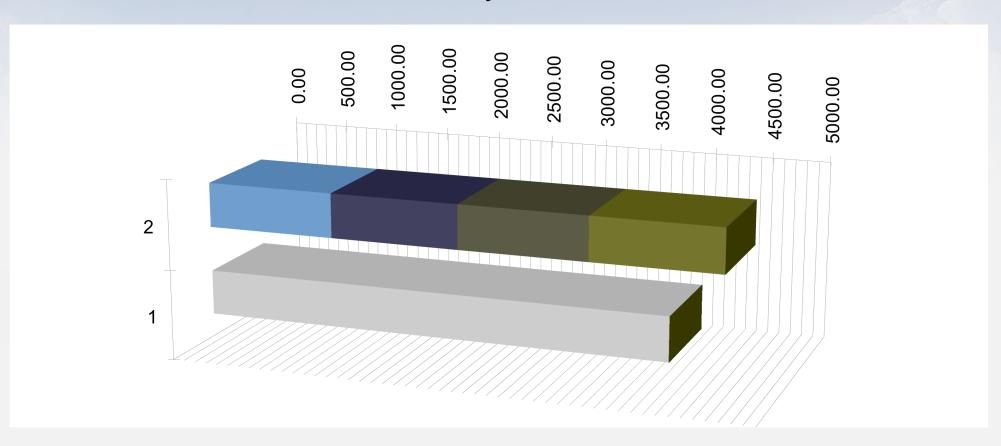




zsmalloc's spillover



bytes



PAGE

Object

■Object

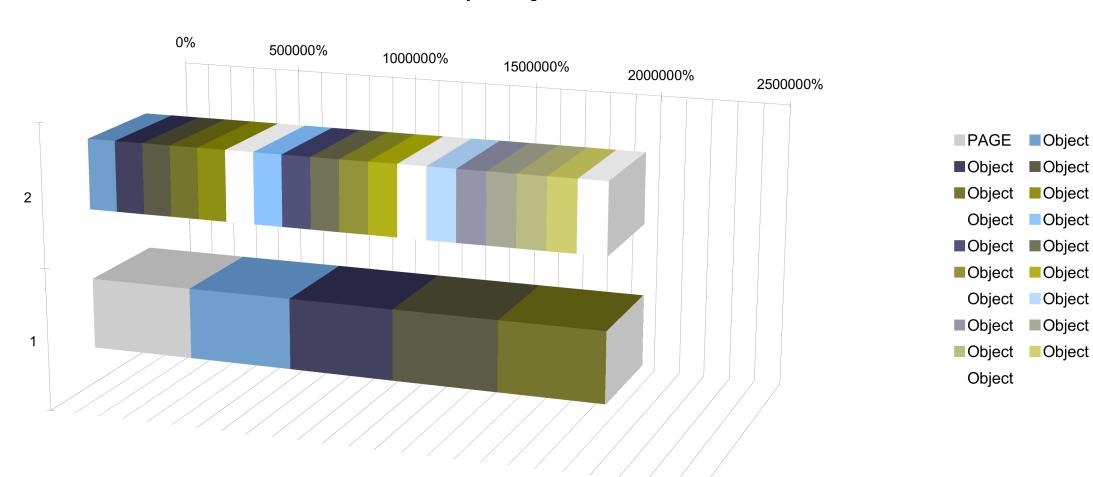
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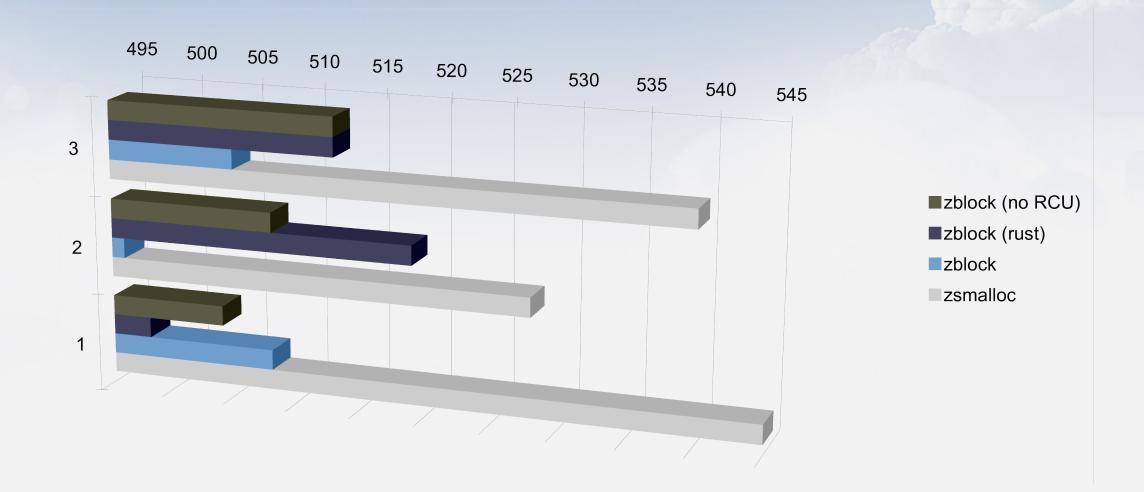
zblock: no spillover







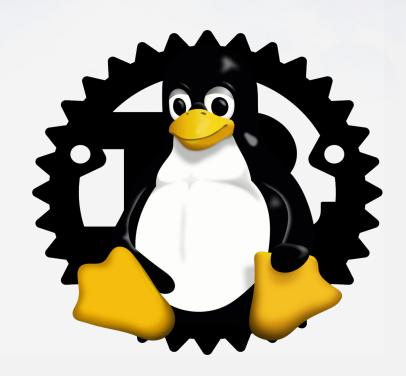
Quick performance comparison Kensulko



Rust allocator: summary

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- Not for all architectures
- New algorithm (page block based)
- Fast and scalable
- Higher memory safety
- Native Rust API
- Implemented, submission under way



Rust compressed ramdisk



- Why
 - ZRAM is bloated
 - Blocks of unused code
- How
 - Zblock.rs provides a safe native Rust API
 - NeoZRAM uses that API
- Status
 - Work ongoing

Shmem and Rust?



- Shmem can't be moved to Rust (should work on all architectures)
- Rust DRM drivers could benefit
- Using shmem for DRM buffer objects is suboptimal
 - BOs often span across very many pages
 - BO pages are treated independently
 - However, BOs are swapped in/out in their entirety
- Idea: implement streaming compression for large BOs
 - No need to use special allocators like zsmalloc
 - Some preliminary work has started

Thanks for your attention!

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