

TOKYO, JAPAN / DECEMBER 11-13, 2025

## SFrame for Arm64 Reliable Stacktrace

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## Agenda

- What is a Stacktrace?
- Why Reliable Stacktrace?
- Providing Reliable Stacktrace
- Current Reliable Stacktrace/Livepatch Support for Arm64
- SFrame as a Solution



# What is a Stacktrace?





#### What is a Stacktrace?

- A stacktrace is a record of functions called by a thread of execution, as recovered from the call stack.
- Typically they appear in a human-readable format along with other diagnostic information

```
[<c12ba080>] ? dump stack+0x44/0x64
[<c103ed6a>] ? warn+0xfa/0x120
[<c109e8a7>] ? module put+0x57/0x70
[<c109e8a7>] ? module put+0x57/0x70
[<c103ee33>] ? warn slowpath null+0x23/0x30
[<c109e8a7>] ? module put+0x57/0x70
[<f80ca4d0>] ? gp8psk fe set frontend+0x460/0x460 [dvb us
[<c109f617>] ? symbol put addr+0x27/0x50
[<f80bc9ca>] ? dvb usb adapter frontend exit+0x3a/0x70 [d
[<f80bb3bf>] ? dvb usb exit+0x2f/0xd0 [dvb usb]
[<c13d03bc>] ? usb disable endpoint+0x7c/0xb0
[<f80bb48a>] ? dvb usb device exit+0x2a/0x50 [dvb usb]
[<c13d2882>] ? usb unbind interface+0x62/0x250
[<c13620d8>] ? __device_release driver+0x78/0x120
[<c1362907>] ? driver detach+0x87/0x90
[<c1361c48>] ? bus remove driver+0x38/0x90
[<c13d1c18>] ? usb deregister+0x58/0xb0
[<c109fbb0>] ? SyS delete module+0x130/0x1f0
[<c1055654>] ? task work run+0x64/0x80
[<c1000fa5>] ? exit to usermode loop+0x85/0x90
[<c10013f0>] ? do fast syscall 32+0x80/0x130
[<c1549f43>] ? sysenter past esp+0x40/0x6a
```



## What is a Stacktrace?

- Conventionally, these traces are made by following the saved frame pointer as a linked-list up the stack
- The saved return address is used to identify the function associated with each frame

Return Address Frame Pointer Local Variables Saved Registers Return Address Frame Pointer **Local Variables** Saved Registers Return Address Frame Pointer Local Variables Saved Registers

Return Address

Frame Pointer

Registers

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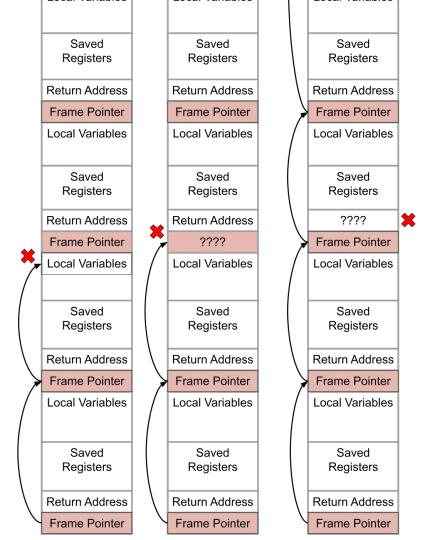
## What is an **Unreliable** Stacktrace?

A stacktrace can be unreliable if it:

- Does not contain the full stack
- Contains incorrect frame/function data
- Contains duplicate entries

How does this happen? More on this later...





# Why Reliable Stacktrace?





## Why Reliable Stacktrace?

Stacktraces are generally useful for profiling and debugging:

- KDump analysis
- Warnings
- Ftrace stack tracing
- Perf event profiling

The benefit of these tools is diminished (but still potentially useful) when the stack itself becomes unreliable.

#### Refs

https://www.man7.org/linux/man-pages/man2/perf\_event\_open.2.html

https://docs.kernel.org/trace/ftrace.html#s tack-trace

https://docs.kernel.org/admin-guide/kdump/kdump.html



## Why Reliable Stacktrace? – Livepatch

Livepatch is a system for applying kernel fixes without rebooting.

- Implemented as a module containing replacement functions
- Control flow is redirected to replacement functions using the dynamic ftrace framework
- Consistency is enforced on a per-task level

Livepatch relies on **reliable stacktraces** to implement this consistency model

#### Refs

https://docs.kernel.org/livepatch/livepatch.html#



## Why Reliable Stacktrace? - The Livepatch Consistency Model

Upon enabling a patch, the Livepatch subsystem enters a transition state

- Individual tasks switch over to the patched state confirmed not to be executing any affected functions
- Transition state only ends when all tasks have been switched to the target patch state

Livepatch applies three distinct strategies to enforce its consistency model

#### Refs

https://docs.kernel.org/livepatch/livepatch.html#consistencv-model



## Why Reliable Stacktrace? – The Livepatch Consistency Model

The first two are partial solutions that **do not** rely on stacktracing:

- 1. Kernel exit switching. A task is switched when it returns to user space from a system call/user space IRQ/signal.
  - Patching I/O-bound user tasks which are sleeping on an affected function. SIGSTOP and SIGCONT will exit the kernel.
  - Patching CPU-bound user tasks the next time they get interrupted by an IRQ.
- 2. Call klp\_update\_patch\_state() on current task directly
  - Useful for idle "swapper" tasks, which never exit the kernel.
     Call in the klp\_update\_patch\_state() in idle loop.
  - No complete solution exists for kthreads.

Refs

https://docs.kernel.org/livepatch/livepatch. html#consistency-model



## Why Reliable Stacktrace? – The Livepatch Consistency Model

#### The third is essential for livepatch success:

- 3. Check the stacks of sleeping tasks. If no patched functions are on a task's stack mark it as patched. Needs HAVE\_RELIABLE\_STACKTRACE
  - This is the most effective strategy to apply patches, and usually succeeds after one try
  - Architectures without HAVE\_RELIABLE\_STACKTRACE
     are not considered fully supported for livepatch.

#### Refs

https://docs.kernel.org/livepatch/livepatch.html#consistency-model



# Providing Reliable Stacktrace





Identifying successful termination of stack trace

- Unwinding ends at an expected kernel entry point
- Location of final stack frame matches expectation
- Unwinder did not exit early with error

#### Refs



Interrupts and exceptions can occur at any point in a function's lifetime.

This can cause inconsistencies on the stack:

- Frame pointer can be invalid during function prologue/epilogue
- Return address may be held in a general purpose register (not on the stack)

#### Refs



Ftrace and kprobe trampoline can result in the rewriting/obscuring of the return address. Thankfully, this problem is solved with generic helpers in both cases:

- ftrace\_graph\_ret\_addr()
- kretprobe find ret addr()

#### Refs



Link register unreliability – the value of the link register may not be updated in lock-step with the setup/teardown of the stack frame.

• Can result in "double counted" functions in the stack trace

#### Refs



## Providing Reliable Stacktrace - The ORC Unwinder

On x86, the ORC unwinder provides reliable stacktraces

- At build time, objtool generates a mapping from instruction address (.org\_unwind\_ip) to stack metadata (.orc\_unwind).
- Creates an off-band method of unwinding the stack

#### Refs

https://docs.kernel.org/arch/x86/orc-unwinder.html



## Providing Reliable Stacktrace - The ORC Unwinder

- Stack unwind is then done reliably using this metadata
- With Frame Pointers disabled completely, this has performance benefits over using frame pointer
- Enables reliable unwind, even across exception boundaries

However, Arm64 is missing the required support from objtool so ORC unwind tables cannot be generated

#### Refs

https://docs.kernel.org/arch/x86/orc-unwinder.html



## Providing Reliable Stacktrace – The ORC Unwinder

A prior RFC proposed to add Arm64 support to objtool to generate ORC tables, and enable reliable stacktrace

- Proposed a "dynamic" form of frame pointer validation
  - The frame pointer on the stack is validated against the frame pointer offset pulled from the ORC table
  - If the calculated FP does not match the expected one, the unwind is considered unreliable
- Involved substantial reorganization of the objtool/ORC code to separate the architecture-generic vs architecture-specific components

#### Refs

https://lore.kernel.org/live-patching/2023 0202074036.507249-1-madvenka@linux.mi crosoft.com/

"Madhavan T. Venkataraman" <<u>madvenka@linux.microsoft.com</u>>



## Providing Reliable Stacktrace - The ORC Unwinder

Feedback from the upstream community indicated a problem with this approach:

- Objtool needs to reverse-engineer the control flow generated by the compiler, creating the possibility that objtool will fail to do so for certain edge cases (e.g. jump tables)
- As code generation strategies evolve, objtool will need to be maintained to keep with the changes, making this approach inherently brittle.

#### Refs

https://lore.kernel.org/live-patching/ZByJ mnc%2FXDcqQwoZ@FVFF77S0Q05N.camb ridge.arm.com/

Mark Rutland < mark.rutland@arm.com>



## Providing Reliable Stacktrace - The ORC Unwinder

A solution using data produced directly from the compiler would be preferred.

- This feedback aligns with past feedback to avoid reverse-engineering control flow: <a href="https://lpc.events/event/11/contributions/971/">https://lpc.events/event/11/contributions/971/</a>
- SFrame was proposed as an alternative solution.

#### Refs

https://lore.kernel.org/live-patching/ZByJ mnc%2FXDcqQwoZ@FVFF77S0Q05N.camb ridge.arm.com/

Mark Rutland < mark.rutland@arm.com>







In v6.12 a patch series landed adds the concept of a "Meta" frame record to explicitly identify key points of interest on the stack:

- Exception boundaries
- Final frames

Marks the two missing pieces to meeting the

HAVE\_RELIABLE\_STACKTRACE requirements.

#### Refs

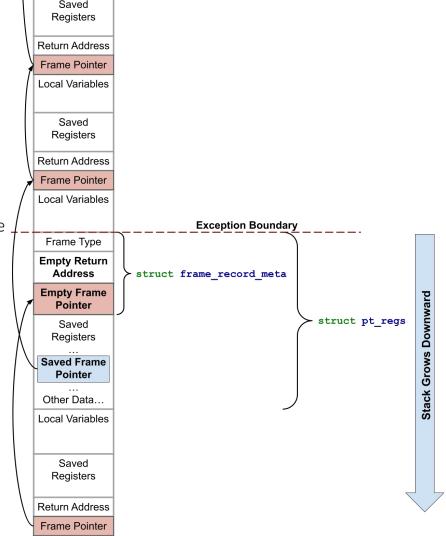
https://lore.kernel.org/linux-arm-kernel/2 0241017092538.1859841-1-mark.rutland@arm.com/

Mark Rutland < mark.rutland@arm.com>



FRAME\_META\_TYPE\_PT\_REGS indicates the presence of a saved struct pt\_regs

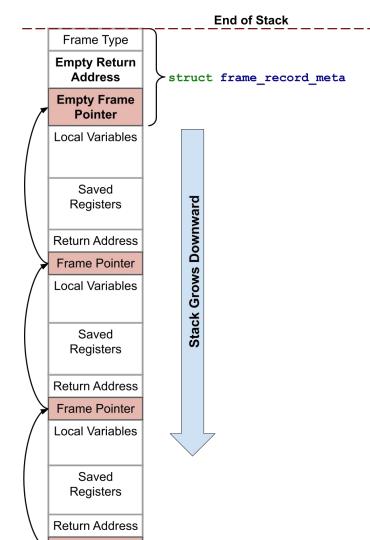
- Occurs at the exception boundary
- Special handling needed to recover the saved Frame
   Pointer and Instruction Pointer





FRAME\_META\_TYPE\_FINAL Indicates the final stack frame

 Once validated to match the task's struct pt\_regs, this confirms successful termination of the stacktrace





Source tracking in the unwind state indicates how each frame was recovered

- "C" from caller of unwind function
- "T" from blocked task's saved PC
- "P" from a pt\_regs::pc

Possible 2nd source

- "F" from fgraph
- "K" from kretprobes

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```

```
Call trace:
show_stack+0x20/0x38 (CF)
dump_stack_lvl+0x60/0x80 (F)
dump stack+0x18/0x28
nmi_cpu_backtrace+0xfc/0x140
nmi_trigger_cpumask_backtrace+0x1c8/0x200
arch_trigger_cpumask_backtrace+0x20/0x40
sysrq_handle_showallcpus+0x24/0x38 (F)
__handle_sysrq+0xa8/0x1b0 (F)
handle_sysrg+0x38/0x50 (F)
pl011_int+0x420/0x570 (F)
__handle_irq_event_percpu+0x60/0x220 (F)
handle_irq_event+0x54/0xc0 (F)
                                         Refs
handle_fasteoi_irg+0xa8/0x1d0 (F
generic_handle_domain_irg+0x34/0x58 (F)
gic_handle_irg+0x54/0x140 (FK)
call_on_irg_stack+0x24/0x58 (F)
                                         https://lore.kernel.org/linux-arm-kernel/2
do_interrupt_handler+0x88/0xa0
                                         0241017092538.1859841-1-mark.rutland@ar
el1_interrupt+0x34/0x68 (F)
el1h_64_irg_handler+0x18/0x28
                                          m.com/
el1h_64_irg+0x6c/0x70
default_idle_call+0x34/0x180 (P)
                                         Mark Rutland < mark.rutland@arm.com >
default_idle_call+0x28/0x180
do_idle+0x204/0x268
cpu_startup_entry+0x3c/0x50 (F)
rest_init+0xe4/0xf0
start_kernel+0x738/0x740
  _primary_switched+0x88/0x98
```

"L" would also indicate when the current frame was recovered from the Link Register, but this feature was removed

- The LR is not always live when an exception occurs
- LR unwinding has to be skipped, meaning unwinding is still unreliable at the exception boundary

#### Refs

https://lore.kernel.org/all/20241211140704.2 498712-2-mark.rutland@arm.com/

Mark Rutland < mark.rutland@arm.com >



A follow-up patch series completes support for reliable stacktrace and livepatch **without sframe** 

- When an exception boundary is detected, stacktrace is marked as unreliable
- Instead of providing/guaranteeing reliable stacktrace in all cases, it
  is sufficient for arch\_stack\_walk\_reliable to report
  when an individual unwind is unreliable
- In practice, a livepatch can be reliably landed with a second unwind pass

#### Refs

https://lore.kernel.org/live-patching/2025 0320171559.3423224-1-song@kernel.org/

Song Liu <song@kernel.org>



This can be considered a partial solution

- Exception boundaries are detected but not yet reliably unwound
- It remains to be seen how reliable this solution is on its own

It should also be noted that this type solution can only help the livepatch use case

 Other applications (and humans) gain no benefit if exceptions are still not reliable!

#### Refs

https://lore.kernel.org/live-patching/2025 0320171559.3423224-1-song@kernel.org/

Song Liu <song@kernel.org>







#### SFrame as a Solution for Reliable Stacktrace

SFrame provides a generalized, toolchain-based solution, inspired by the ORC unwinder

- Provides an assembler-generated .sframe section, which contains a two-level lookup structure
- Similar to ORC, but avoids reverse-engineering the binary

#### Refs

https://sourceware.org/binutils/docs/sfra me-spec.html

https://static.sched.com/hosted\_files/ossn a2023/d1/0SS\_LinuxCon\_SFrame\_final.pdf? gl=1

Indu Bhagat <indu.bhagat@oracle.com>



Provides the minimal necessary information to provide a stacktrace:

- Frame Pointer (FP)
- Return Address (RA)
- Canonical Frame Address (CFA)

#### Refs

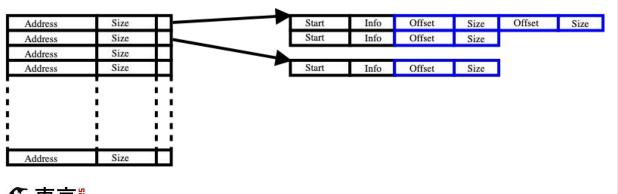
https://sourceware.org/binutils/docs/sfra me-spec.html

https://static.sched.com/hosted\_files/ossn a2023/d1/OSS\_LinuxCon\_SFrame\_final.pdf? gl=1

Indu Bhagat < indu.bhagat@oracle.com >



In full, the SFrame table contains the information necessary to calculate the next frame's register values



Refs

https://lwn.net/Articles/1029189/



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Sframe Function Descriptor Entry (FDE) table provides function metadata:

- Start address
- Size
- FRE count
- Other metadata...

FDEs can be sorted on PC for faster lookup

#### Refs

https://sourceware.org/binutils/docs/sframe-spec.html



Sframe Frame Row Entry (FRE) table is per-FDE and provides a complete record of stack trace information for all instructions in each function.

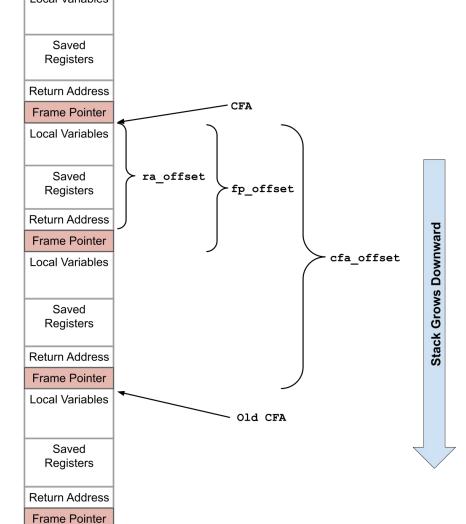
- FRE word info provides generic metadata:
  - CFA BASE\_REG ID (SP or FP)
  - o Other data...
- FRE definitions for Aarch64:
  - CFA = BASE\_REG + offset1
  - o RA = CFA + offset2
  - o FP = CFA + offset3

#### Refs

https://sourceware.org/binutils/docs/sframe-spec.html



With the SFrame information, the locations of the Frame Pointer and Return Address can be calculated reliably at each unwind step.





## Pseudocode unwind algorithm:

```
fre = sframe_find_fre (pc);
if (fre)
    base_reg_val = sframe_fre_base_reg_fp_p (fre) ? fp : sp;
    cfa_offset = sframe_fre_get_cfa_offset (fre);
    ra_offset = sframe_fre_get_ra_offset (fre);
    fp_offset = sframe_fre_get_fp_offset (fre);
    cfa = base_reg_val + cfa_offset;
    next_frame->sp = cfa;
    ra_stack_loc = cfa + ra_offset;
    next_frame->pc = read_value (ra_stack_loc);
    if (fp_offset is VALID)
        fp_stack_loc = cfa + fp_offset;
        next_frame->fp = read_value (fp_stack_loc);
    else
        next_frame->fp = fp;
else
    ret = ERR_NO_SFRAME_FRE;
```

#### Refs

https://sourceware.org/binutils/docs/sframe-spec.html



## SFrame as a Solution – Our Proposal

How to unwind across exception boundaries

- Similar to ORC unwinder:
  - Annotate where the base register is and its relative CFA offset
  - Compiler will generate . sframe FRE provides a CFA offset relative to the base register

#### Refs

https://lore.kernel.org/lkml/20250904223 850.884188-1-dylanbhatch@google.com/

Dylan Hatch < <a href="mailto:dylanbhatch@google.com">dylanbhatch@google.com</a>>
Weinan Liu < <a href="mailto:wnliu@google.com">wnliu@google.com</a>>



- Sframe can support x86\_64, AArch64, s390x
  - We can reuse same sframe unwind functions for different architectures
  - Generated by compiler directly so we don't have to run objtool after binaries are built

Refs

https://sourceware.org/binutils/wiki/sframe



- Complement with FP based approach. Fallback to FP approach if...
  - o 00T kernel module is built without sframe information.
  - SFrame unwind encounters an error and is deemed unreliable

#### Refs

https://lore.kernel.org/lkml/20250904223 850.884188-1-dylanbhatch@google.com/

Dylan Hatch <<u>dylanbhatch@google.com</u>> Weinan Liu <<u>wnliu@google.com</u>>



SFrame for userspace call-stack unwinding support

- Add userspace unwinding support by accessing the SFrame data in the ELF section of the executable from the kernel.
- https://lore.kernel.org/lkml/cover.1730150953.git.jpoimboe@kernel.org/

#### Refs

https://lore.kernel.org/all/20251119132323.1 281768-1-iremus@linux.ibm.com/

https://lore.kernel.org/all/cover.1737511963.git.jpoimboe@kernel.org/

https://lore.kernel.org/all/2025082720154 8.448472904@kernel.org/

Jens Remus < iremus@linux.ibm.com >
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Steven Rostedt < rostedt@kernel.org >



## SFrame as a Solution – Ongoing Improvements

However, SFrame is currently only supported for GCC

- Support for LLVM is an ongoing effort
- Ongoing discussions point to an uncertainty about whether the tradeoffs SFrame presents align with the maintenance burden in LLVM

Toolchain support is essential for broader adoption of SFrame

#### Refs

https://discourse.llvm.org/t/rfc-adding-sfr ame-support-to-llvm/86900



## SFrame as a Solution – Ongoing Improvements

A number of improvements are anticipated for SFrame V3

- Flexible support for topmost frames
- Mark outermost frames

#### Refs

https://sourceware.org/binutils/wiki/sframe/sframev3todo





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