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From Fragmentation to Integration: Enhancing sched_ext BPF Scheduler Interoperability with Linux

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sched_ext User Questions

What workloads is <scheduler> good for?

Is <scheduler> is cache aware?

How does <scheduler> compare to eevdf/cfs?

What is the best scheduler for my hardware?

What is the best scheduler for gaming?



Are there answers?

scx_layered - per workload scheduling policy

general purpose vs specialized (scx_tickless)

What does the user care about (throughput vs latency)?

Do they know what they care about?

What is the system topology?



Is Fragmentation A Problem?

Knowledge silos:

BPF scheduling ≠ kernel scheduling

Duplicated effort:

Each scheduler reimplements topology, idle selection, etc.

Missed opportunities:

Kernel improvements don't automatically benefit BPF schedulers

Innovation barrier:

Hard to contribute ideas across the boundary



Fragmentation

BPF Scheduler Developers:

- Expertise in specific scheduler (scx_lavd, scx_layered, etc.)
- Custom topology representations (~1k LOC in scx_utils)
- Scheduler-specific patterns and idioms

Kernel Scheduler Developers:

- Deep knowledge of sched_domain, load balancing, topology
- Understanding of decades of optimization and edge cases
- Limited visibility into BPF scheduler innovations



The Cost

New developers learn scheduler-specific quirks, not scheduling fundamentals

Improvements in scx_layered don't help scx_rusty (and vice versa)

BPF scheduler expertise doesn't translate to kernel scheduler contribution

Each new scheduler starts from scratch -> good for simple topology



Goal

Bridge these worlds without stifling innovation

Encourage novel scheduling policies and experiments

But build on shared foundations and common vocabulary

Make BPF scheduler development a path to kernel scheduler understanding



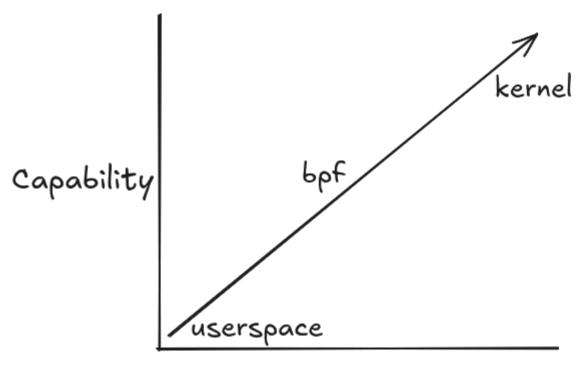
Example: Topology

Problem: Hybrid BPF/userspace schedulers need topology in multiple places

BPF scheduler developers may never see kernel/sched/topology.c

Migration cost/load balancing equally important to scheduling policies in complex topologies







Iteration Time

Userspace Topology

```
pub struct Topology {
  pub nodes: BTreeMap<usize, Node>,
  pub span: Cpumask,
  pub smt enabled: bool,
  // Skip indices for fast access
  pub all Ilcs: BTreeMap<usize, Llc>,
  pub all cores: BTreeMap<usize, Core>,
  pub all cpus: BTreeMap<usize, Cpu>,
```



Can we make it better?

```
// Kernel per-CPU variables exposed via _ksym
extern int sd_llc_size _ksym _weak; // CPUs in LLC
extern struct sched_domain* sd_llc _ksym; // LLC domain ptr
extern struct cpufreq_policy* cpufreq_cpu_data __ksym;
## Helper macros generate accessors
DEFINE_PER_CPU_VAL_FUNC(cpv_llc_id, int, sd_llc_id)
DEFINE_PER_CPU_VAL_FUNC(cpu_llc_size, int, sd_llc_size)
DEFINE_THIS_CPU_VAL_FUNC(cpu_llc_id) // Current CPU
```



Data structure reuse

```
// scx_beerland/src/bpf/main.bpf.c:248
static bool is_cpu_same_llc(s32 this_cpu, s32 that_cpu) {
  return cpu_llc_id(this_cpu) == cpu_llc_id(that_cpu);
}
```

Benefit:

Direct access to kernel topology cache, no userspace required

Bridge to kernel:

Using these accessors exposes developers to kernel topology concepts



What can we do?

Composability/BPF Arenas
Emil (next talk)

More kfuncs/default implementations

Idle CPU selection

Scheduler Governors- Steven Rostedt

Seems like a good long term approach



Al as a Knowledge Bridge

Developer:

"How does the kernel maintain LLC topology?"

AI:

Shows kernel/sched/topology.c, sd_llc per-CPU cache
Shows update_top_cache_domain() and when it's called
Links to sched_domain structure definition

Result:

Hours of exploration condensed to minutes





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