

Linux Plumbers Conference

Vienna, Austria | September 18-20, 2024

Address Space Isolation

x86 Microconference Brendan Jackman <<u>jackmanb@google.com</u>>

Context links: <u>2024 RFC</u> | <u>Demo branch</u> (with tests + optimisations) | <u>LSF/MM/BPF</u> + <u>recordings</u>

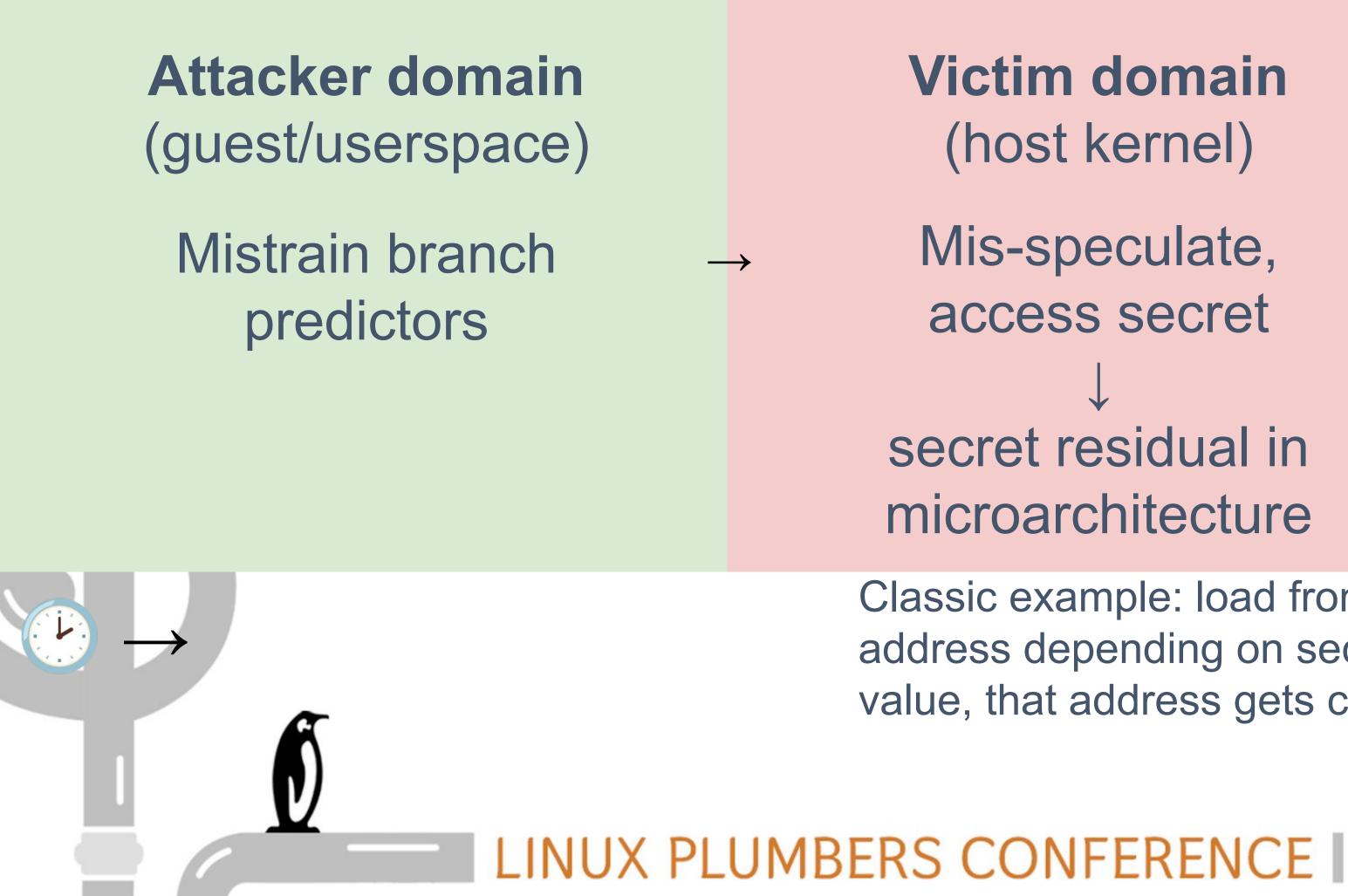
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Agenda

- Hasty background refresher if needed? (5 mins)
- A look at some perf data (2 mins)
- Discuss how to get this thing merged

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(Some) CPU Exploits refresher



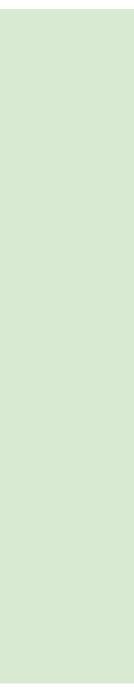
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-speculate, ess secret	
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Classic example: load from address depending on secret value, that address gets cached.

Attacker domain (guest/userspace)

recover secret

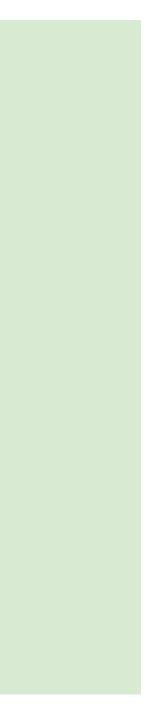
Check access timing to see which address is cached. "Flush+Reload"





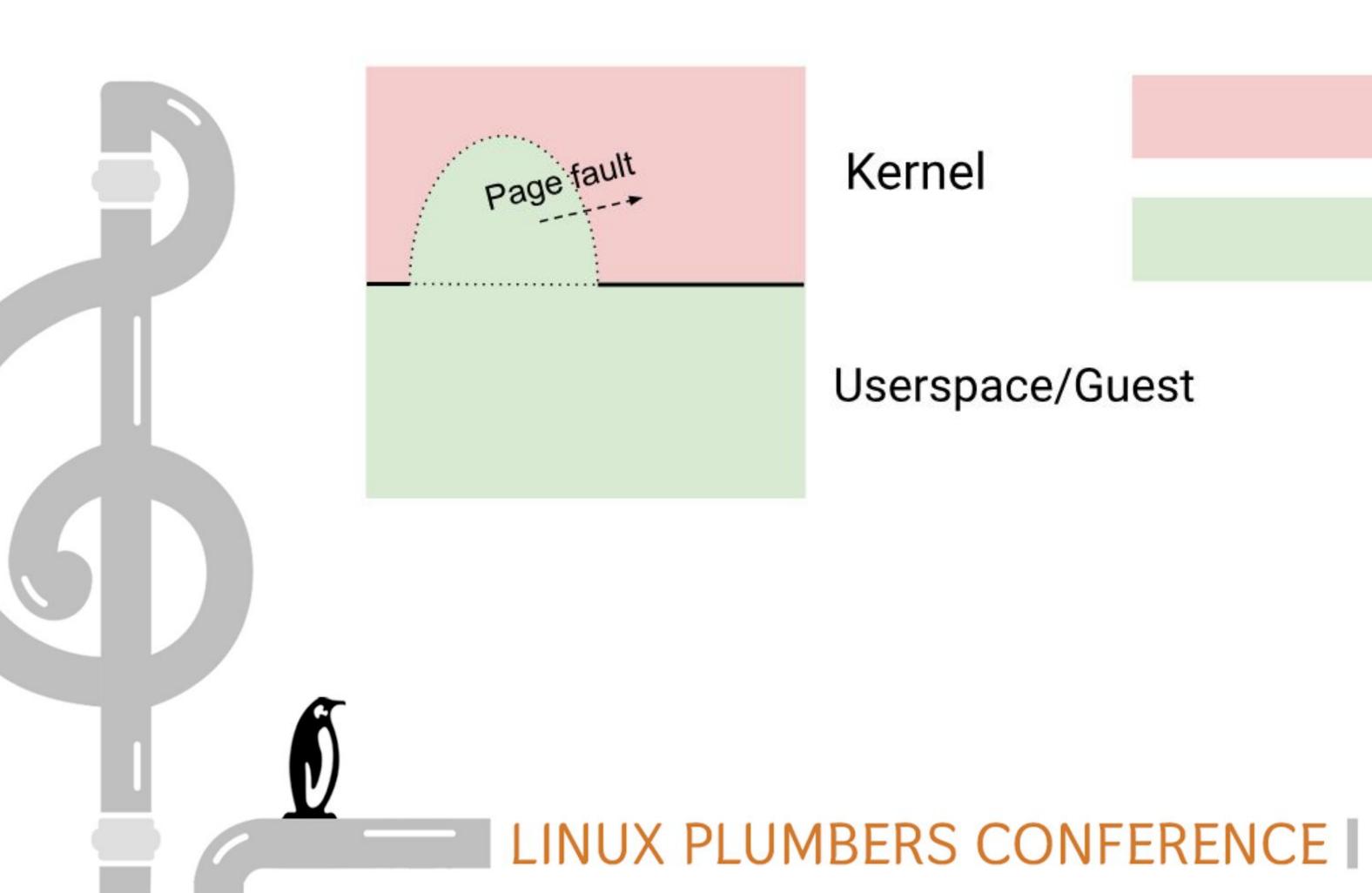


(Some) CPU Exploits refresher



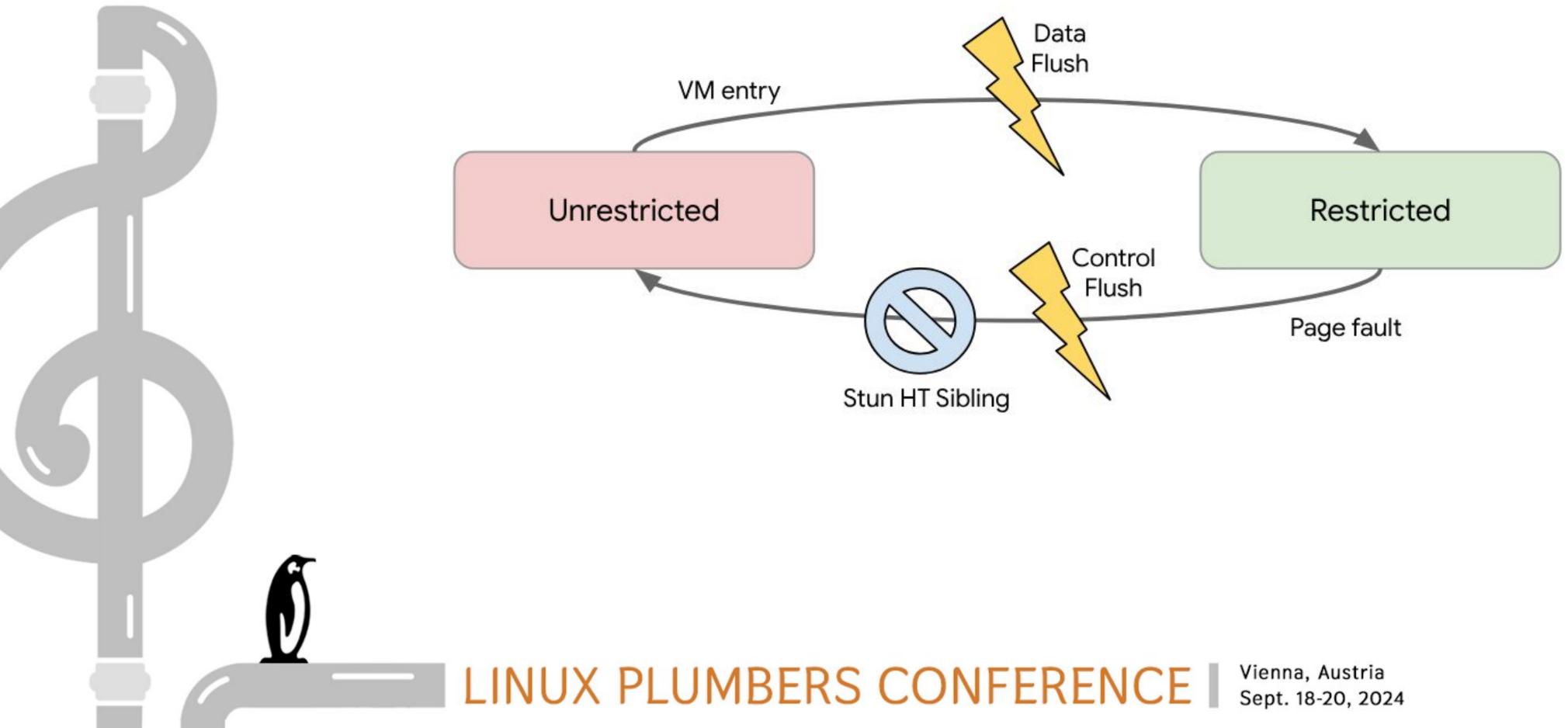


ASI refresher



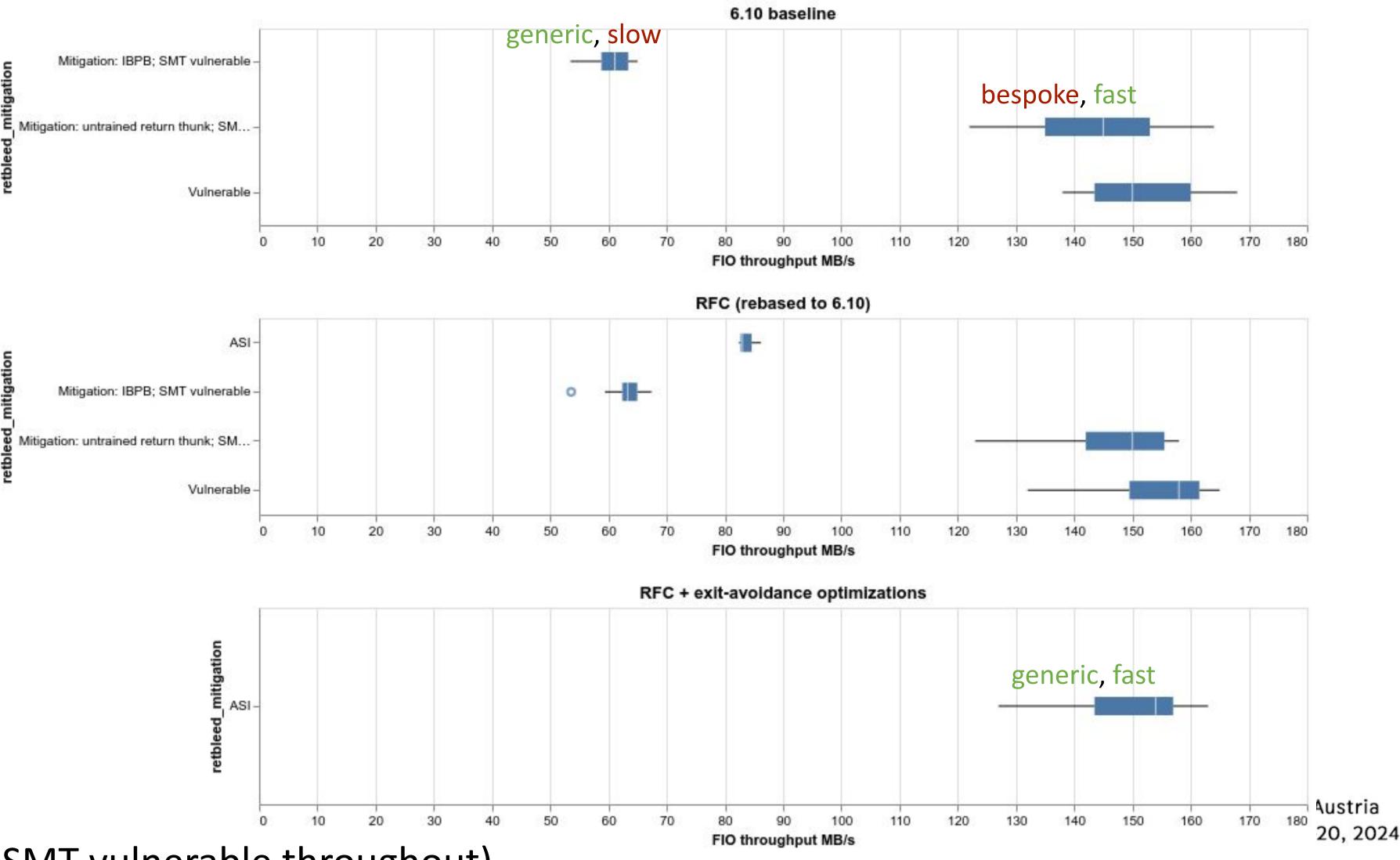


ASI refresher





Performance (Zen2)



(SMT vulnerable throughout)

Performance = comparable to bespoke mitigations

Security properties = comparable to sledgehammer mitigations

I presented something similar at LSF/MM/BPF. Claimed low-confidence in the data. But now it's a different benchmark, different platform, more evidence for same conclusion. Also matches experience in Google's kernel.



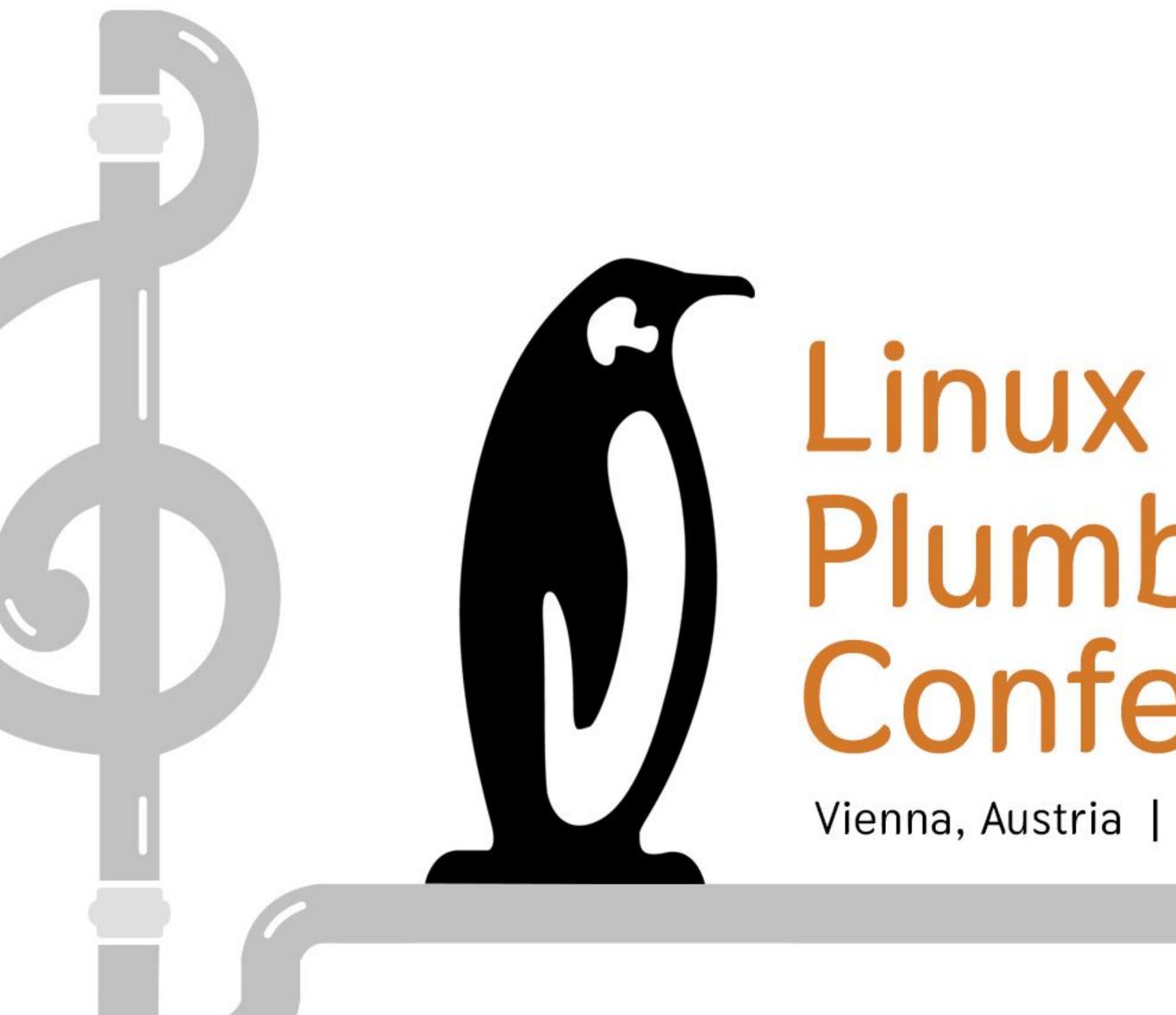
Topics for discussion



- mm folks seemed up for it at LSF/MM/BPF (but Mel wasn't there)
- "Denylist" approach: start with only protecting GFP USER directmap
 - Probably prevents all existing attacks, but obviously not watertight
- Roadmap for bare-metal sandboxing
- How should users configure it?
- We have tests (KUnit, e2e exploits that stop working)
 - Testing mitigations is hard though

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- But lets us work in-tree on an ASI that's actually viable for production - Build up security from there, with a meaningful performance baseline



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