

# Leverage Homa: Enhancing Homa Linux for Efficient RPC Transportation

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# Agenda

- Homa Introduction
- Limitation of Homa in RPC context
- Homa Congestion Control Enhancements
- Homa RPC Streaming Enhancements
- Future Improvements
- Conclusion
- Q&A



# Homa Introduction



# Why TCP is Wrong for Data Center?

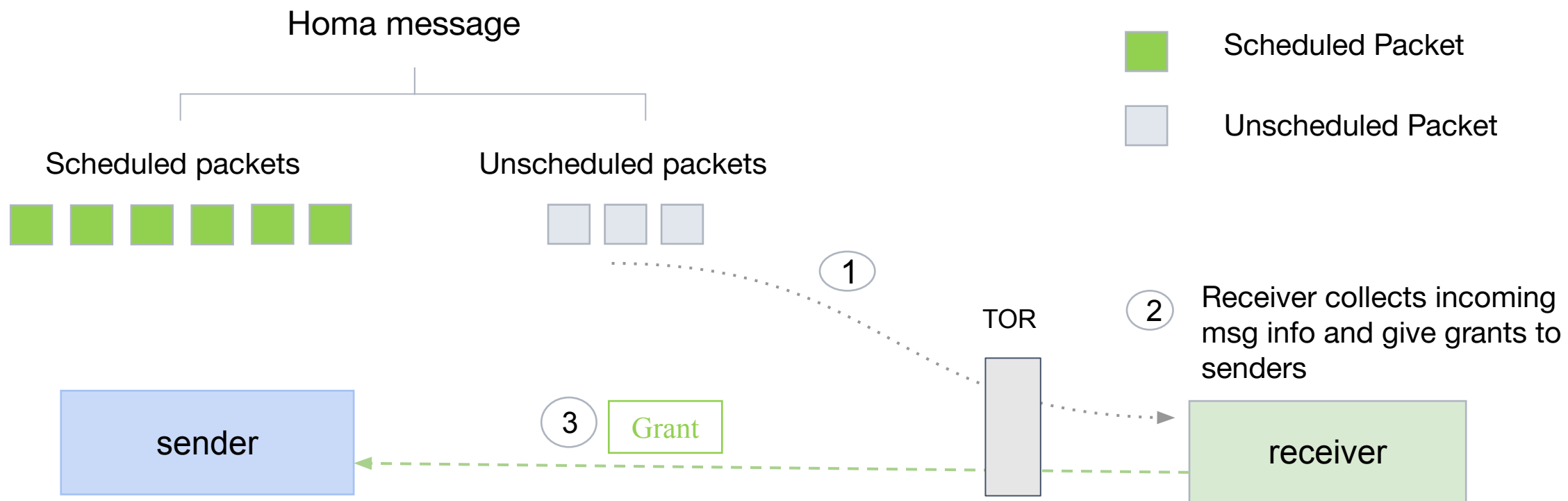
- Designed for wide-area networks
- Connections
- Stream orientation
- Fair scheduling
- Sender-driven congestion control
- It doesn't take advantage of in-network priority queues.
- In-order delivery, restricting opportunities for load balancing



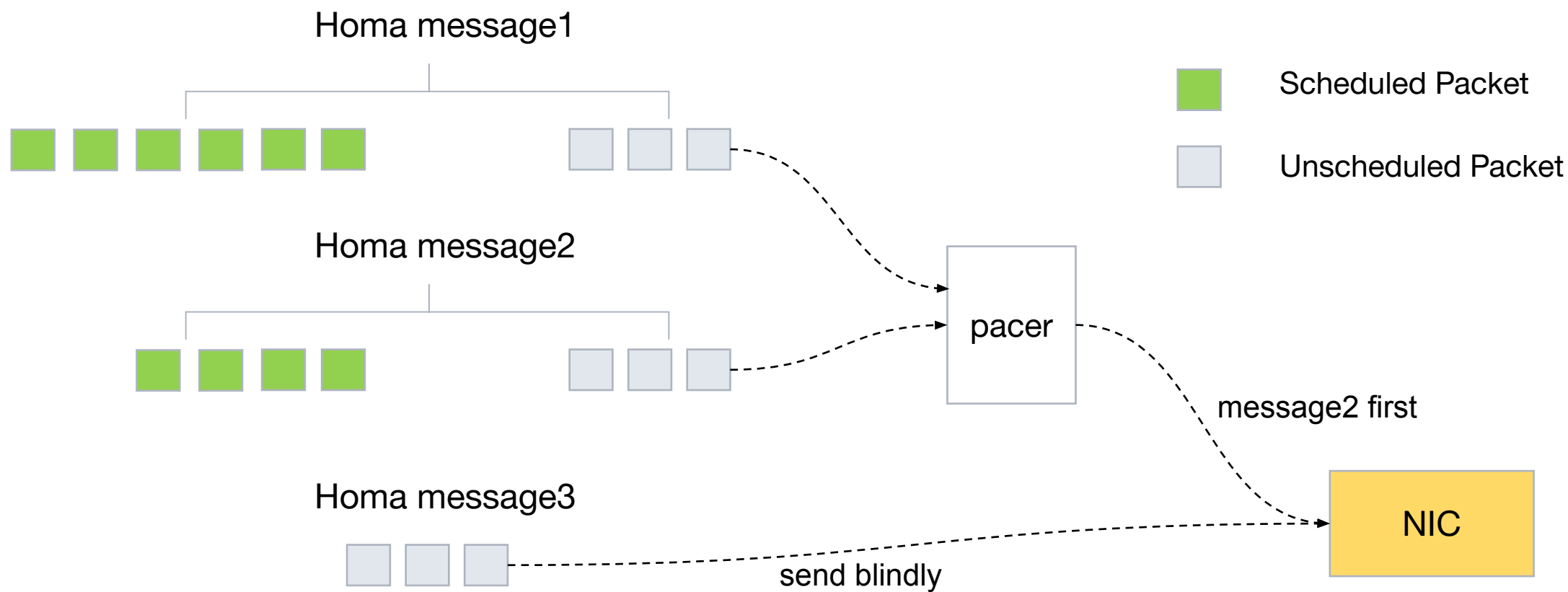
# Homa Protocol Introduction

- Designed for datacenter networks with extremely low latencies
- Connectionless, no connection cost, no long life connection state
- Message based protocol
- SRPT(Shortest Remaining Process Time first)
- Use in-network priorities
- Receiver-driven packet scheduling
- Overcommitment

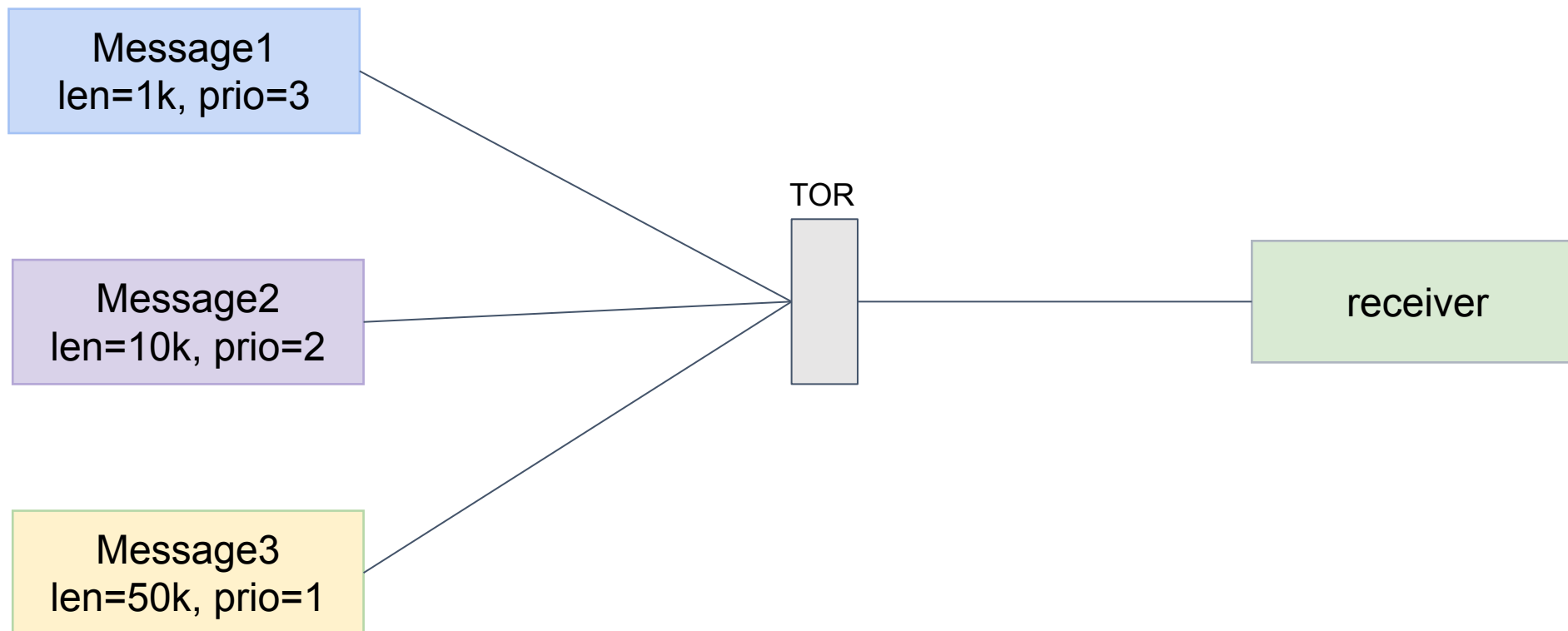
# Overview



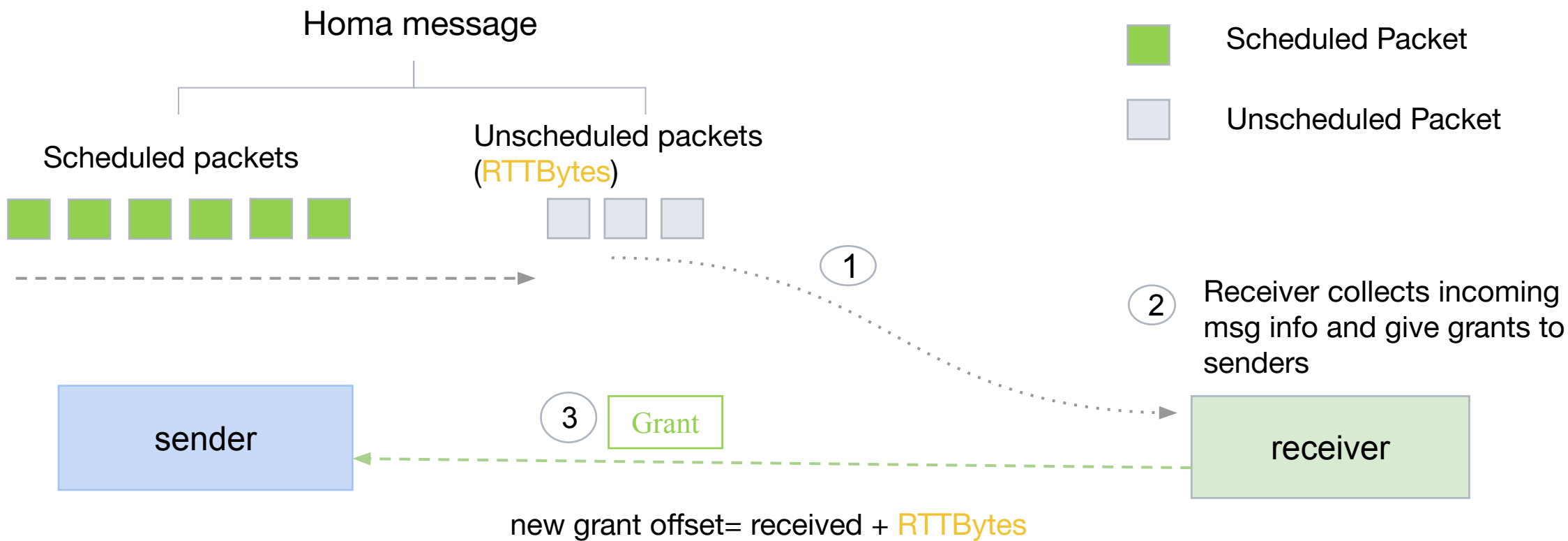
# Sender – SRPT and Pacer



# In-network priorities

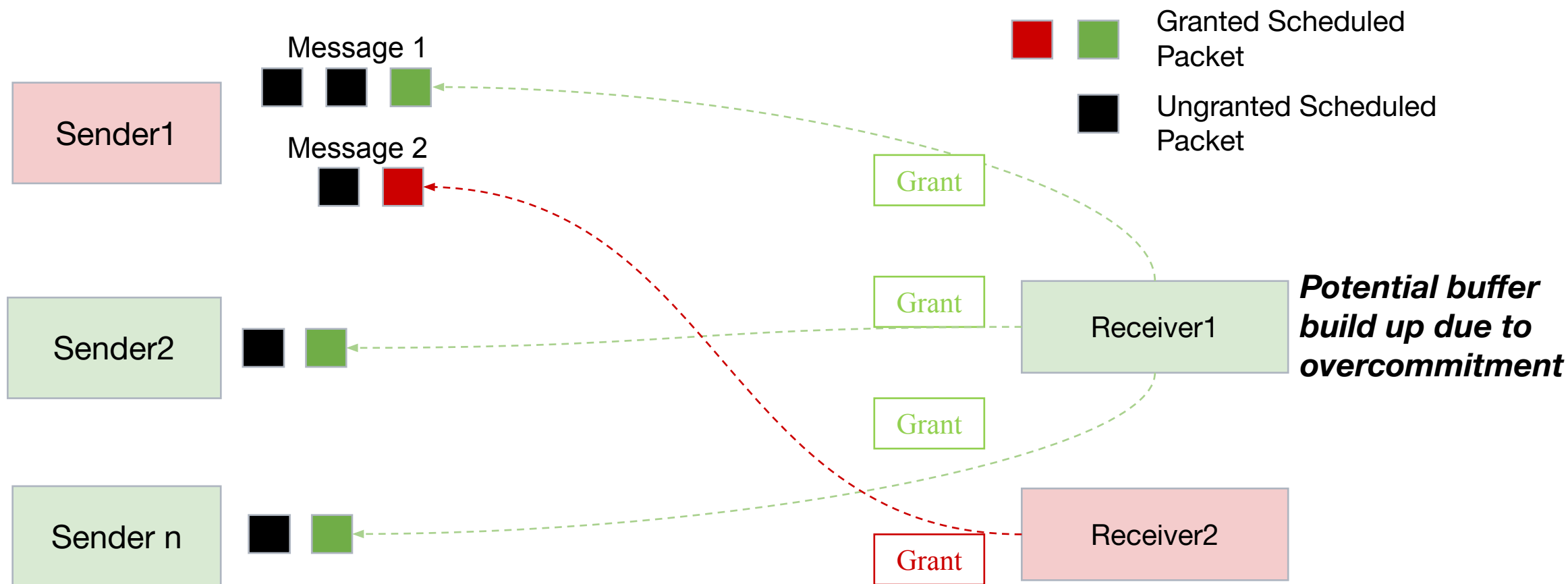


# Server – Packet Scheduling



**$RTTBytes$  is a pre-set fixed value**

# Server – Controlled Overcommitment





# Limitations of Homa as RPC transport protocol



# Limitations of Homa as RPC transport protocol

## Inefficient Pipelining for Large Message

- Homa is message-based protocol, while ensuring complete message delivery, it hinders efficient pipelining. As a result, for large RPC messages(size > 50k), Homa is not as good as TCP.

## Non-standard Socket API interface

- It is not easy to map Homa RPC ID to existing RPC framework.
- No long lived RPC: A stream RPC is consisted of many Homa RPCs, which incurs the overhead of creating and reclaiming them.



# Limitations of Homa as RPC transport protocol

## Performance is sensitive to RTTBytes Config

- Manually config RTTBytes can be inconvenient
- Single preset value is not enough for diverse RTT and receiver downlink bandwidth.

## Weak Congestion Control when RTT is larger than 20 us

- High RTTBytes can more easily lead to incast congestion
- Low RTTBytes is not able to cover RTT



# Limitations of Homa as RPC transport protocol

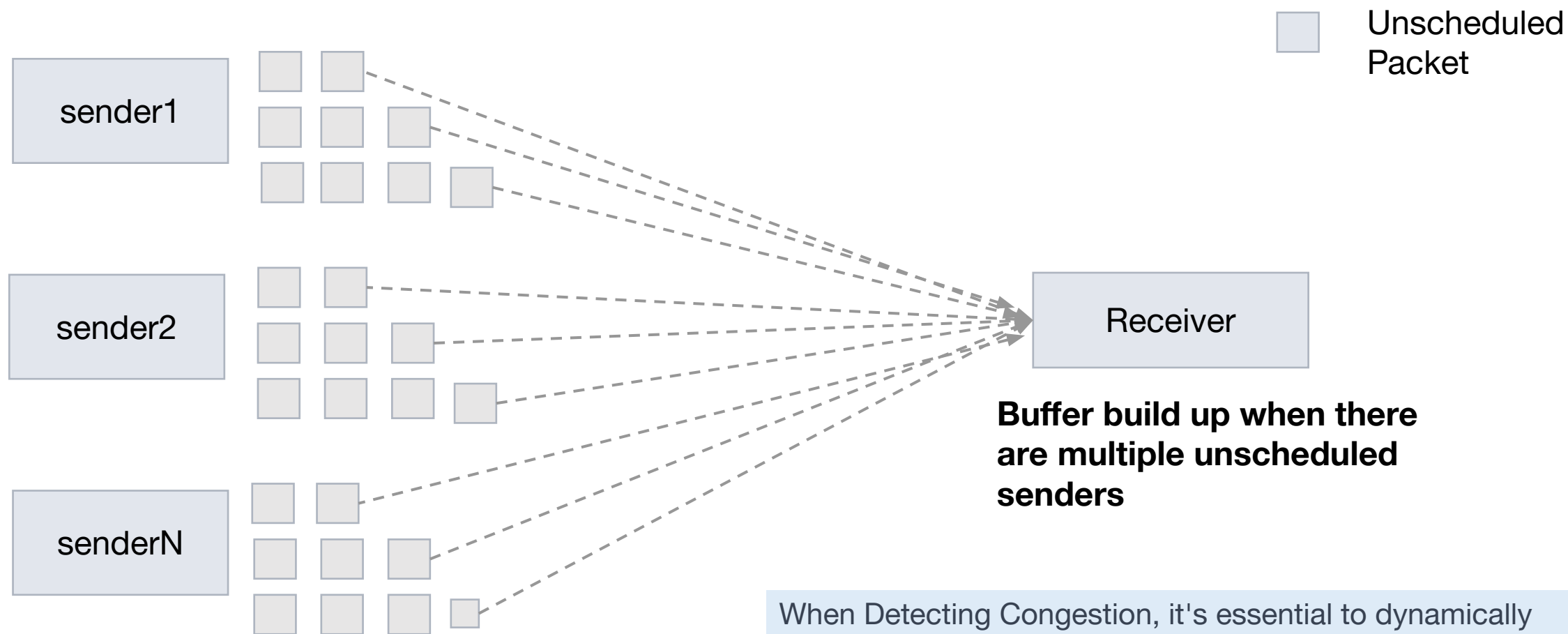
## Homa cannot coexist harmoniously with TCP

- In practice, network resources need to be shared among protocols like TCP
- Homa assumes the bandwidth is all used by itself. If the bandwidth is shared by TCP, Homa maybe over-generous on unscheduled bytes, and overgrant scheduled bytes.

## Homa inactively handle incast

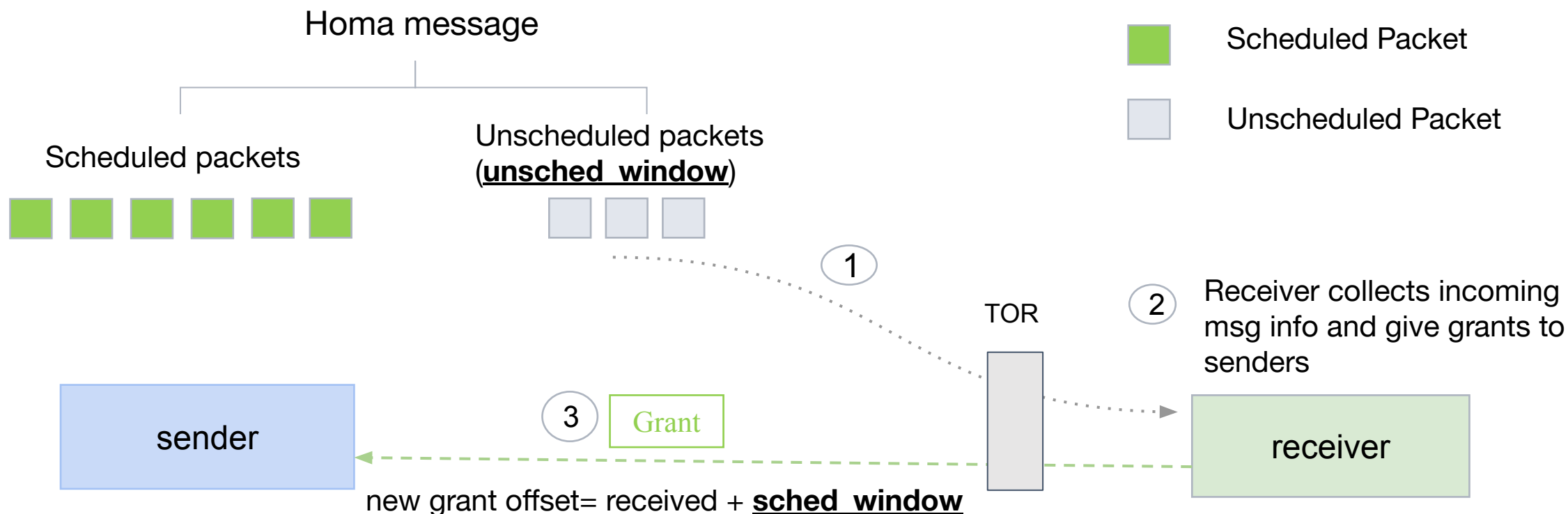
- Homa assumes that the most severe forms of incast are predictable because they are self-inflicted by outgoing RPCs.
- Homa assumes the incast where several machines simultaneously send requests to one server is rare.

# Unscheduled Packet Incasting



When Detecting Congestion, it's essential to dynamically reduce unscheduled Window to prevent congestion

# Static Congestion Window is Insufficient



***Unscheduled Window and Scheduled Window need different values!***



# Homa Congestion Control Enhancements

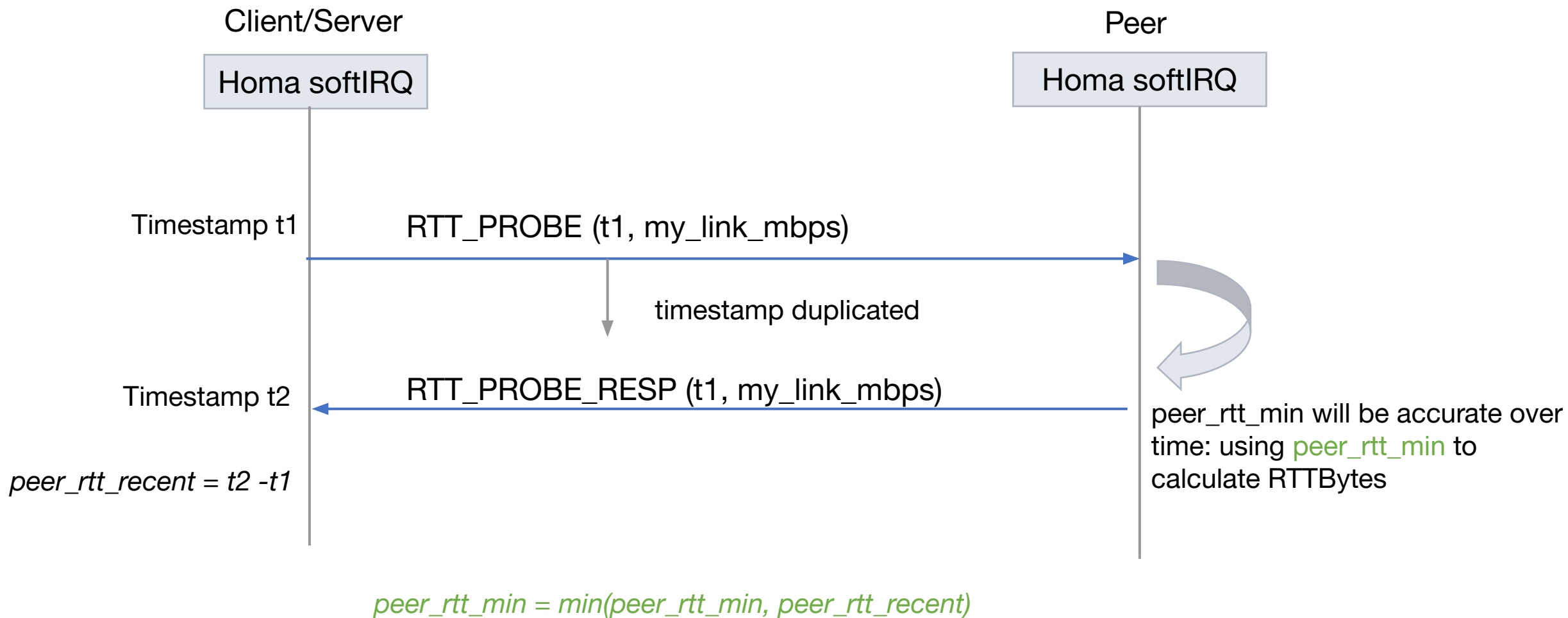


# Homa Congestion Control Enhancements

## Dynamic Per Peer Adjustable Window

- Real-time peer RTT detection
- RTT-informed congestion detection
- Adaptive per-peer adjustable unscheduled/scheduled window based on congestion

# Real-time Peer RTT Detection





# RTT Informed Congestion Detection

peer\_rtt\_min: The minimum RTT value detected over time for this peer

peer\_rtt\_low: The low threshold of RTT

peer\_rtt\_mid: The middle point of RTT

peer\_rtt\_high: The high threshold of RTT

```
peer_rtt_low = peer_rtt_min * 2  
peer_rtt_mid = peer_rtt_low * 2  
peer_rtt_high = peer_rtt_mid * 2  
IF peer_rtt_recent > peer_rtt_high  
    Set congestion To TRUE  
ENDIF
```

# Sender - Dynamic adjusting unscheduled window

```
unsched_window = peer_rtt_recent * peer_link_mbps / 8
```

# Sender - Dynamic adjusting unscheduled window

~~$unsched\_window = peer\_rtt\_recent * peer\_link\_mbps / 8$~~

IF  $peer\_rtt\_recent < peer\_rtt\_mid$

$peer\_rtt\_unsched = peer\_rtt\_low$

ELIF  $peer\_rtt\_recent < peer\_rtt\_high$

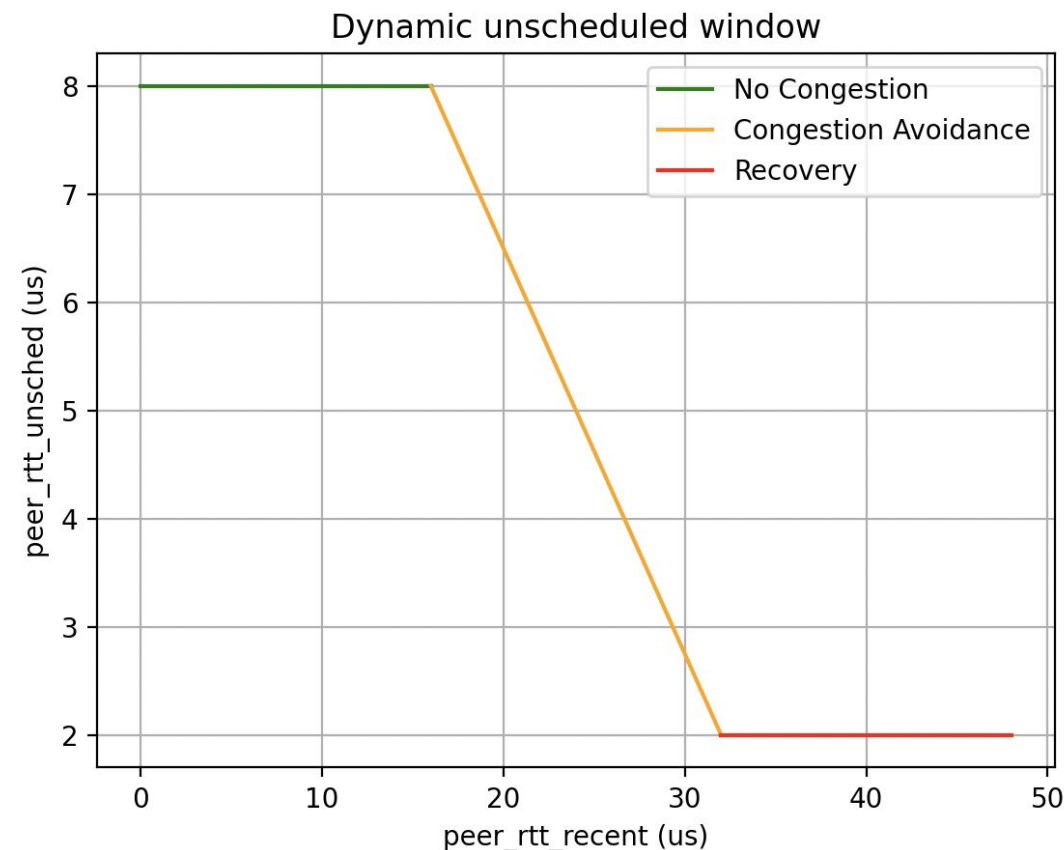
$peer\_rtt\_unsched = peer\_rtt\_low -$   
 $(peer\_rtt\_recent - peer\_rtt\_mid) * 3 / 8$

ELSE

$peer\_rtt\_unsched = rtt\_min / 2$

ENDIF

$unsched\_window = peer\_rtt\_unsched * peer\_link\_mbps / 8$



# Sender - Unscheduled Ratio

*Unscheduled\_ratio = Total unscheduled packets / Total length of all messages*

*IF `unscheduled\_ratio <= 40%` AND `pacers throttle list is not empty`*

*SET unsched\_window To unsched\_window / 2*

*ENDIF*

# Receiver - Scheduled Window

Use *peer\_rtt\_grant* as reference RTT to calculate scheduled window for grant:

```
peer_rtt_grant = 3 * peer_rtt_min;
```

```
scheduled_window = peer_rtt_grant * local_link_mbps / 8;
```

```
max_incomming = scheduled_window * max_overcommit;
```



# Homa Congestion Control Enhancements

## Performance is nearly independent to RTTBytes Config

- Instead of fixed and static RTT, we now use real-time RTT.

## Homa harmoniously coexist with TCP

- Homa can be aware of the existence of other protocols' traffic by feeling the turbulence of RTT.

## Homa actively handle incast

- Incast can be reflected by high RTT, then senders can adjust the unscheduled window dynamically.



# Performance Evaluation



# Testbed Setup

## 25G network

CPU: Intel(R) Xeon(R) Platinum 8163 (96 core, 2.50GHz)

RAM: 400G DIMM DDR4

NIC: Mellanox ConnectX-4 Lx 25 Gbp

TOR Switch; Arista DCS-7050SX3-48YC12-F 25G ports

## 100G network

CPU: Intel(R) Xeon(R) Silver 4314 (64 cores, 2.4 GHz)

RAM: 400 GB DIMM DDR4

NIC: Mellanox Technologies MT28841 dual-port 100Gb/s

TOR Switch: Ruijie Networks RG-S6580-48CQ8QC 100G ports

# Testbed Setup

Name	Mean	Description
W2	433	Search application at Google [33].
W3	2423	Aggregated workload from all applications running in a Google datacenter [33].
W4	60175	Hadoop cluster at Facebook [32].
W5	385315	Web search workload used for DCTCP [1].

We use the same workload in [2] to test. Since Homa congestion control enhancements only have a trivial effect when message size is small, we focus on workload W4, W5 and other fixed-size long messages.

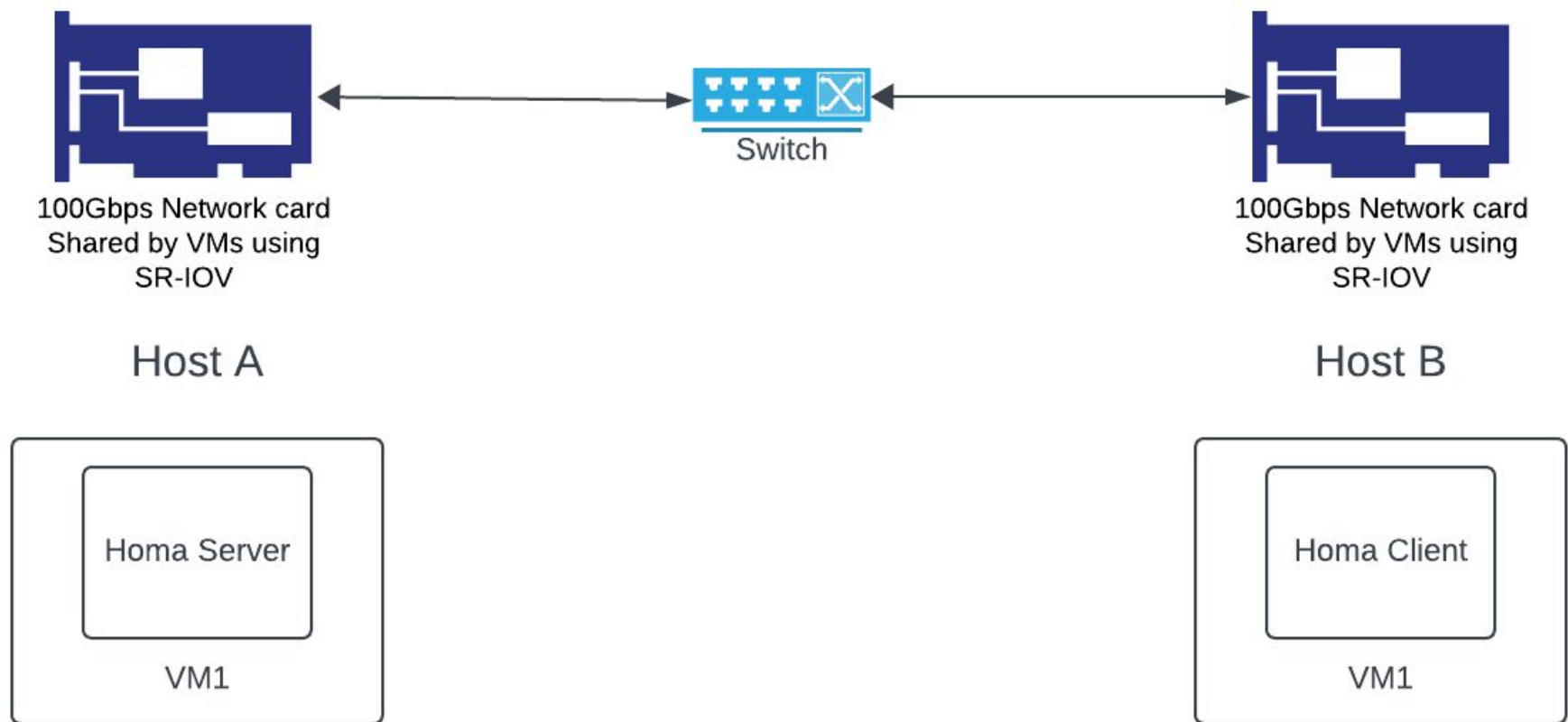


# Test Tool

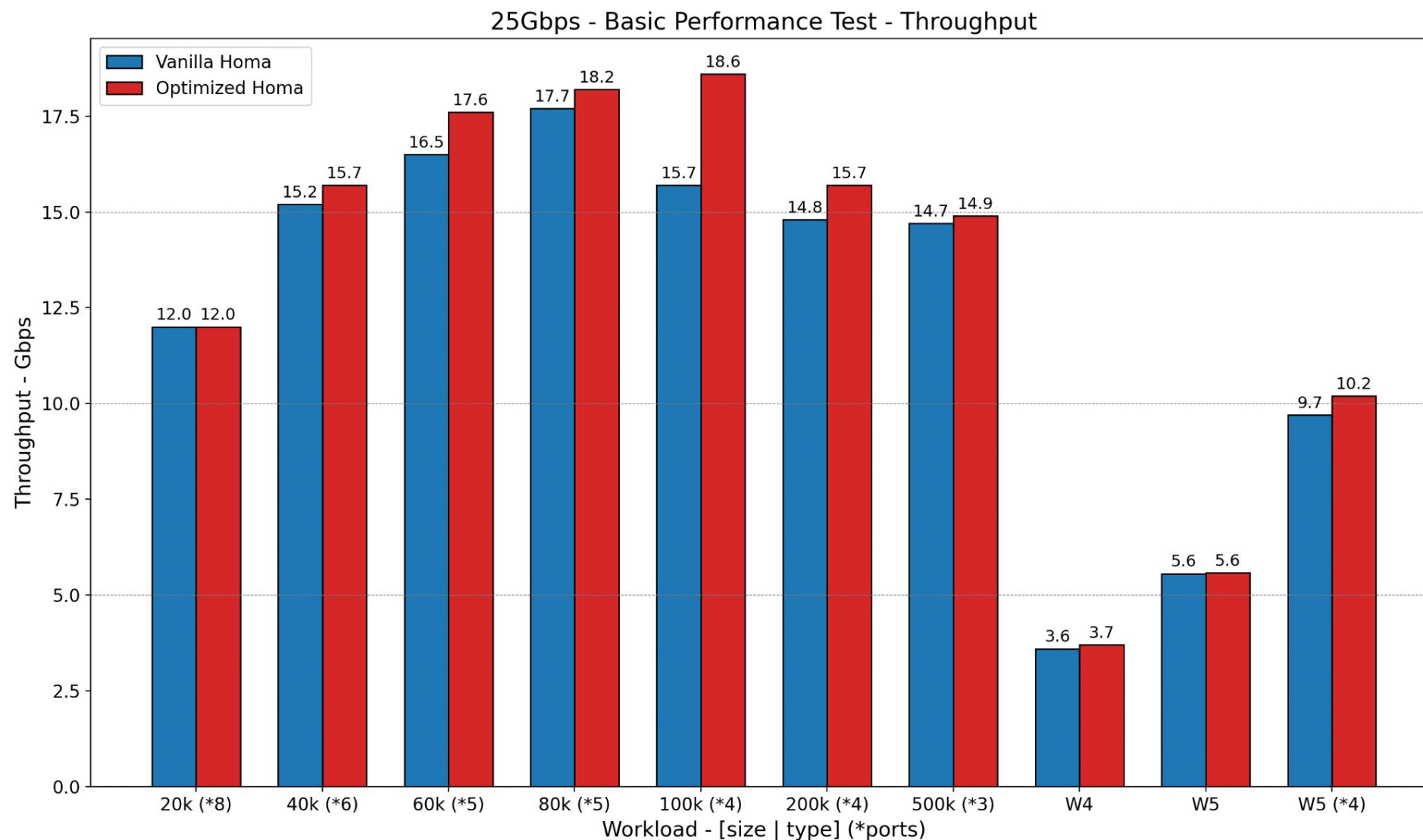
Test application [cp\\_node](#) is a program to test the performance(including throughput, latency, etc) of Homa or TCP. In our test, we mainly tweak some parameters for clients to adjust the behavior of the client node.

- workload, workload to run the test, could be fixed-size or workload type.
- ports, for clients, the number of ports on which to send requests.

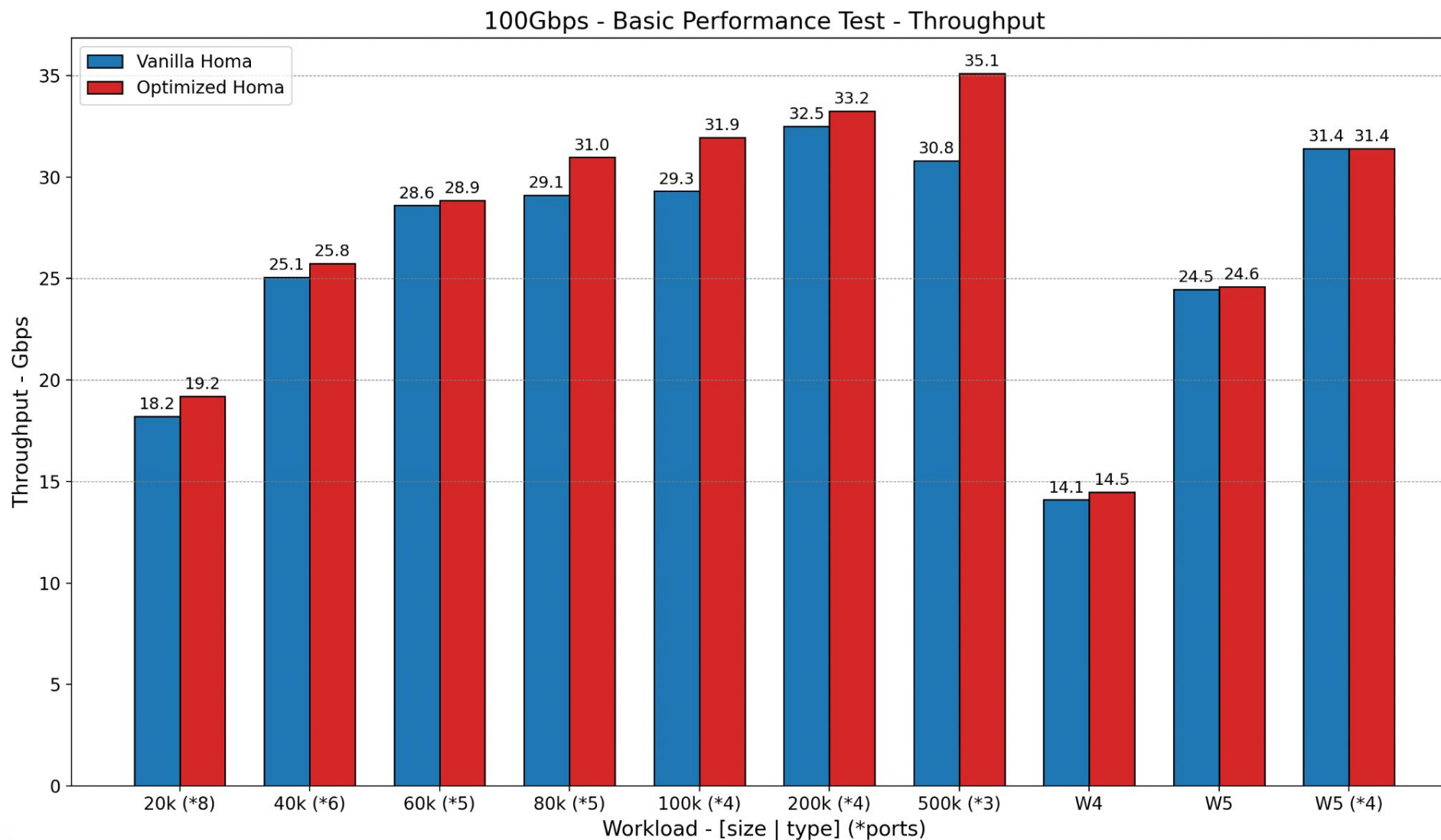
# Basic Performance Setup



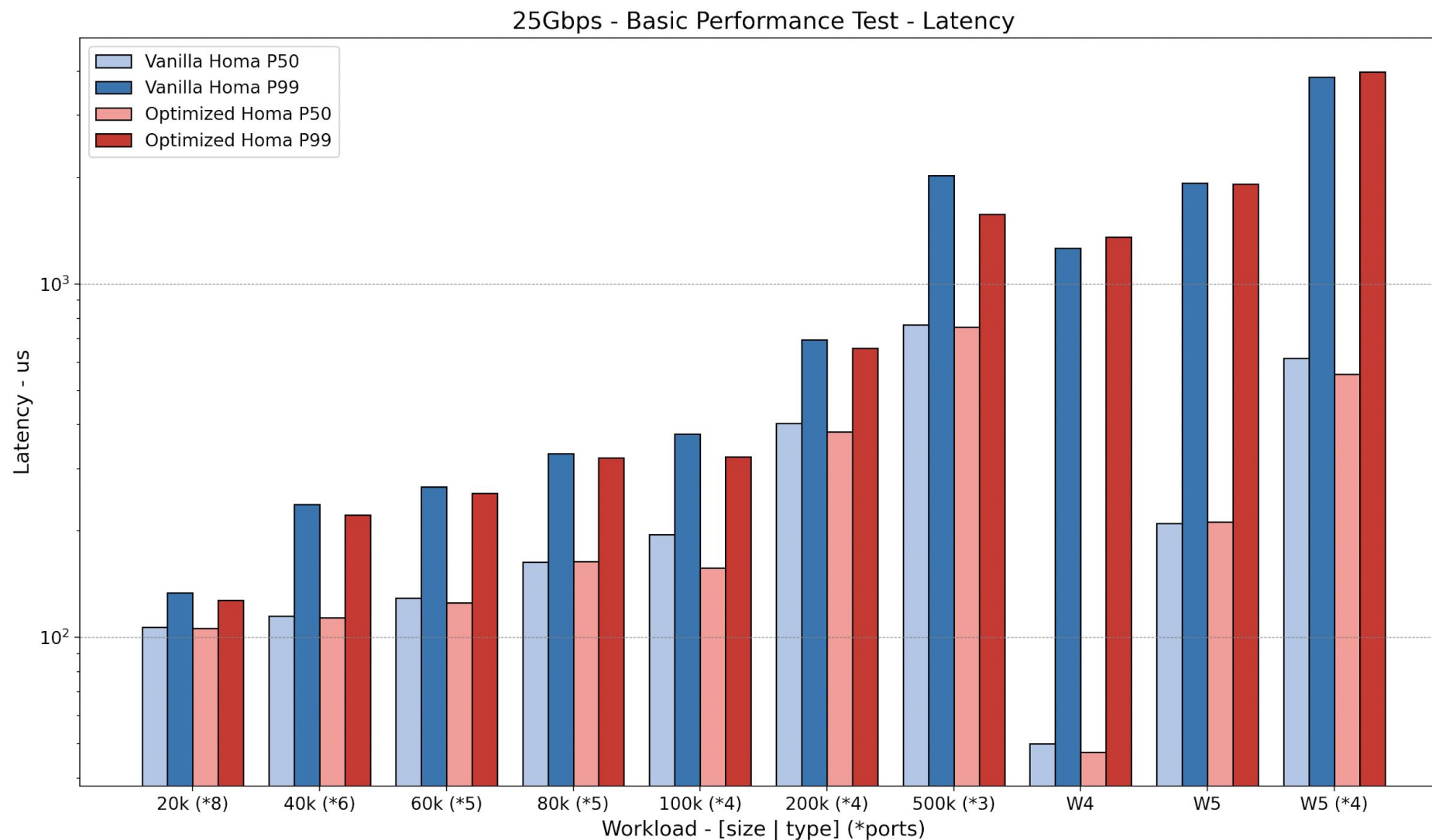
# Basic Performance - 25Gbps - Throughput



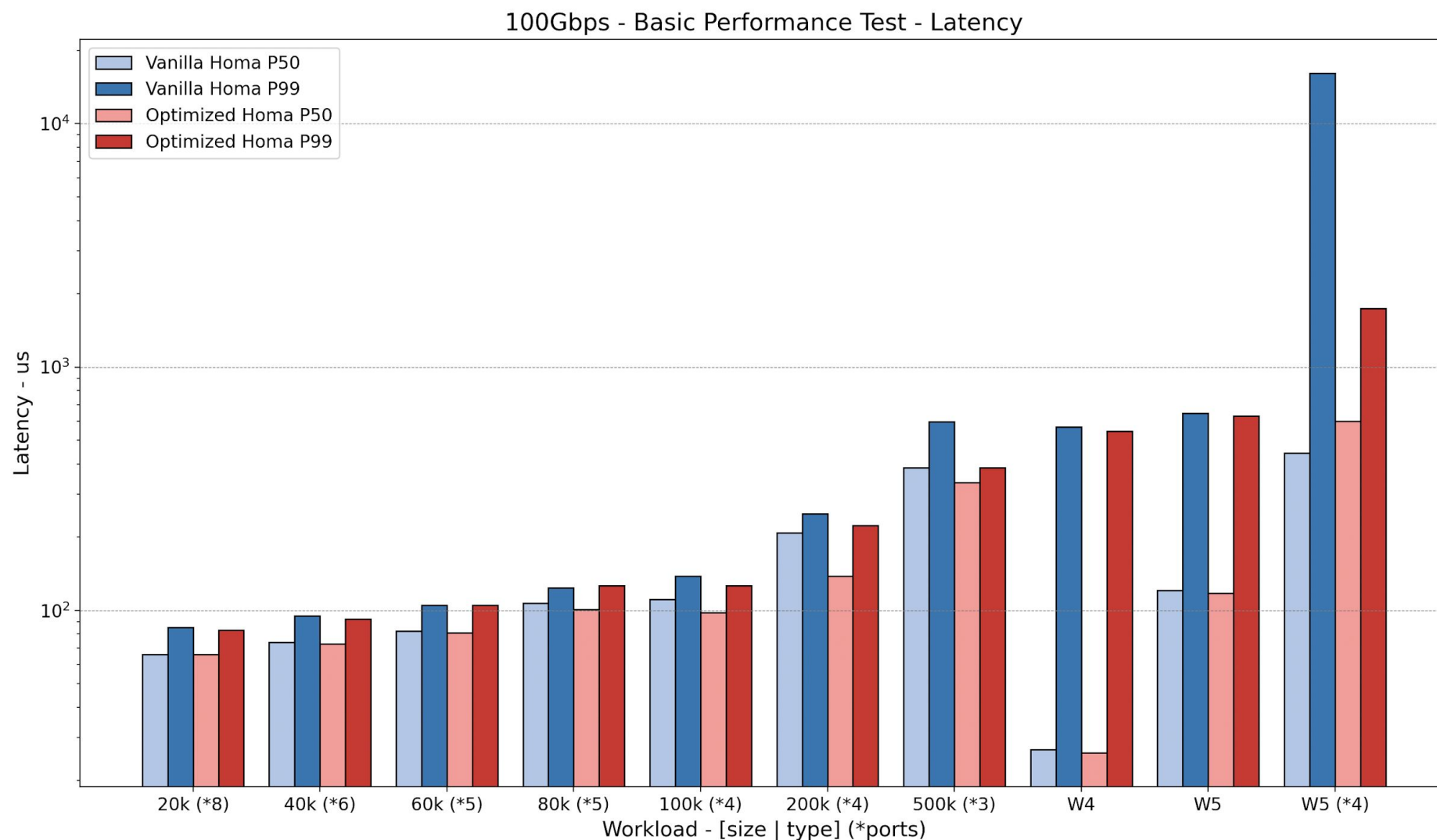
# Basic Performance - 100Gbps - Throughput



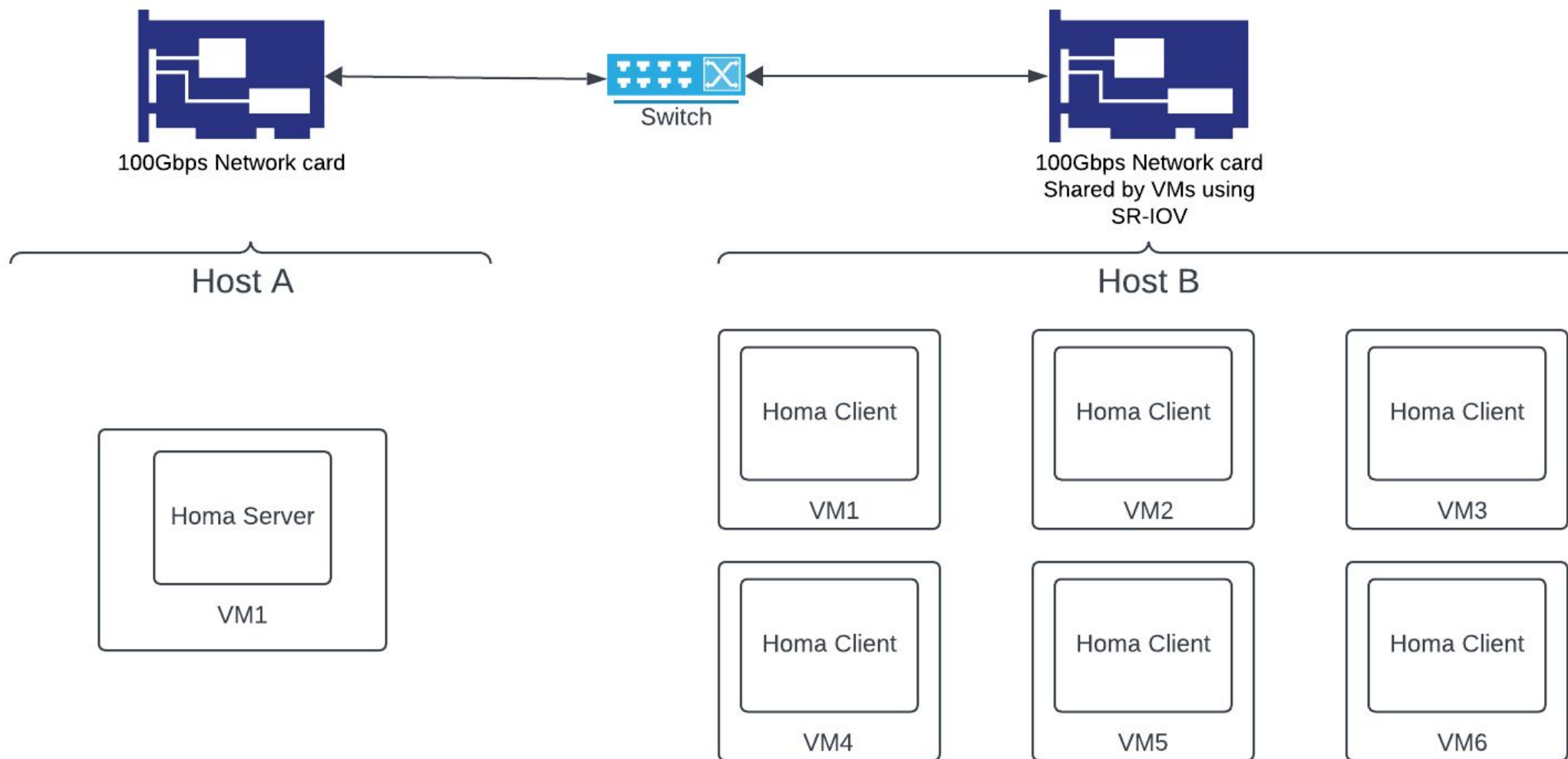
# Basic Performance - 25Gbps - Latency



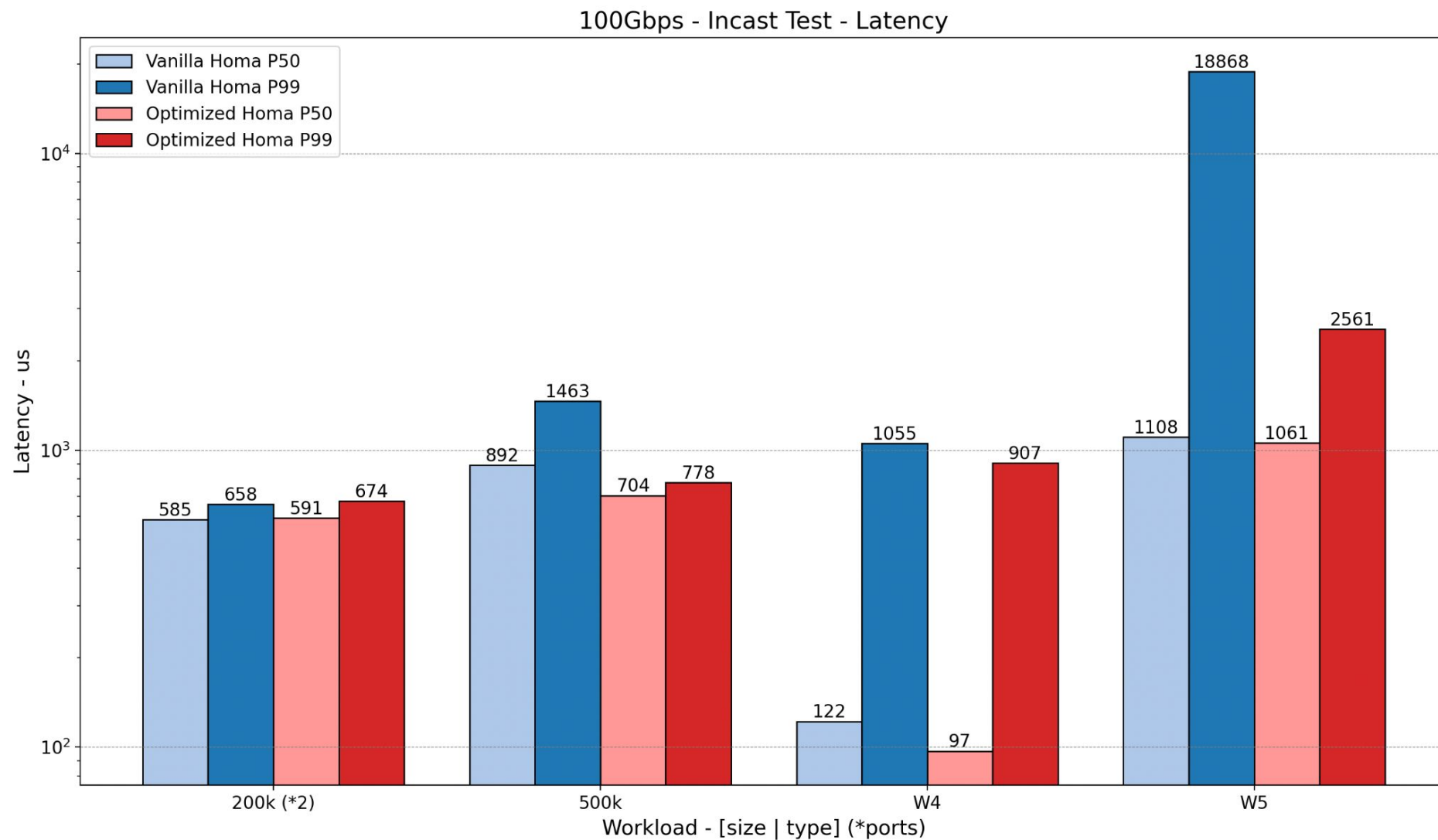
# Basic Performance - 100Gbps - Latency



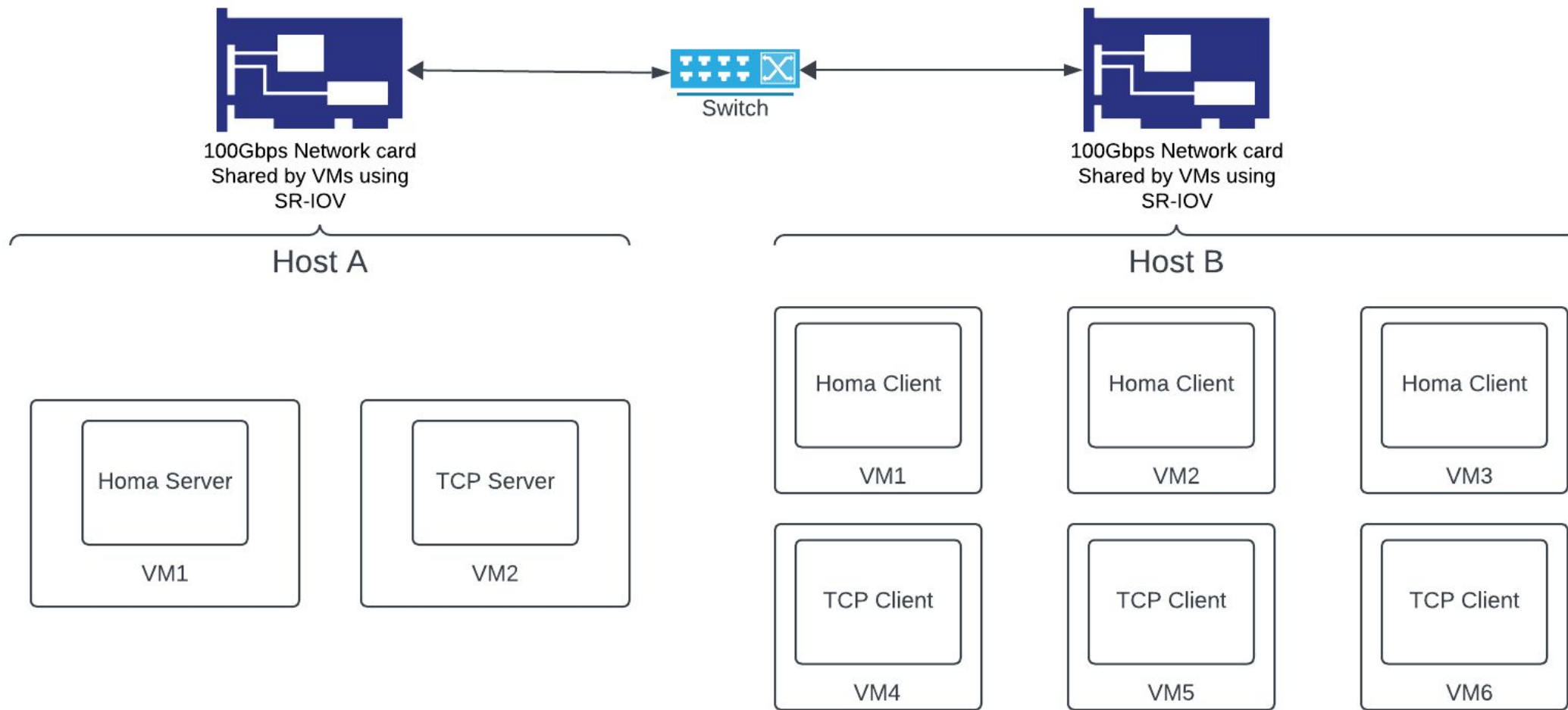
# Incast - Setup



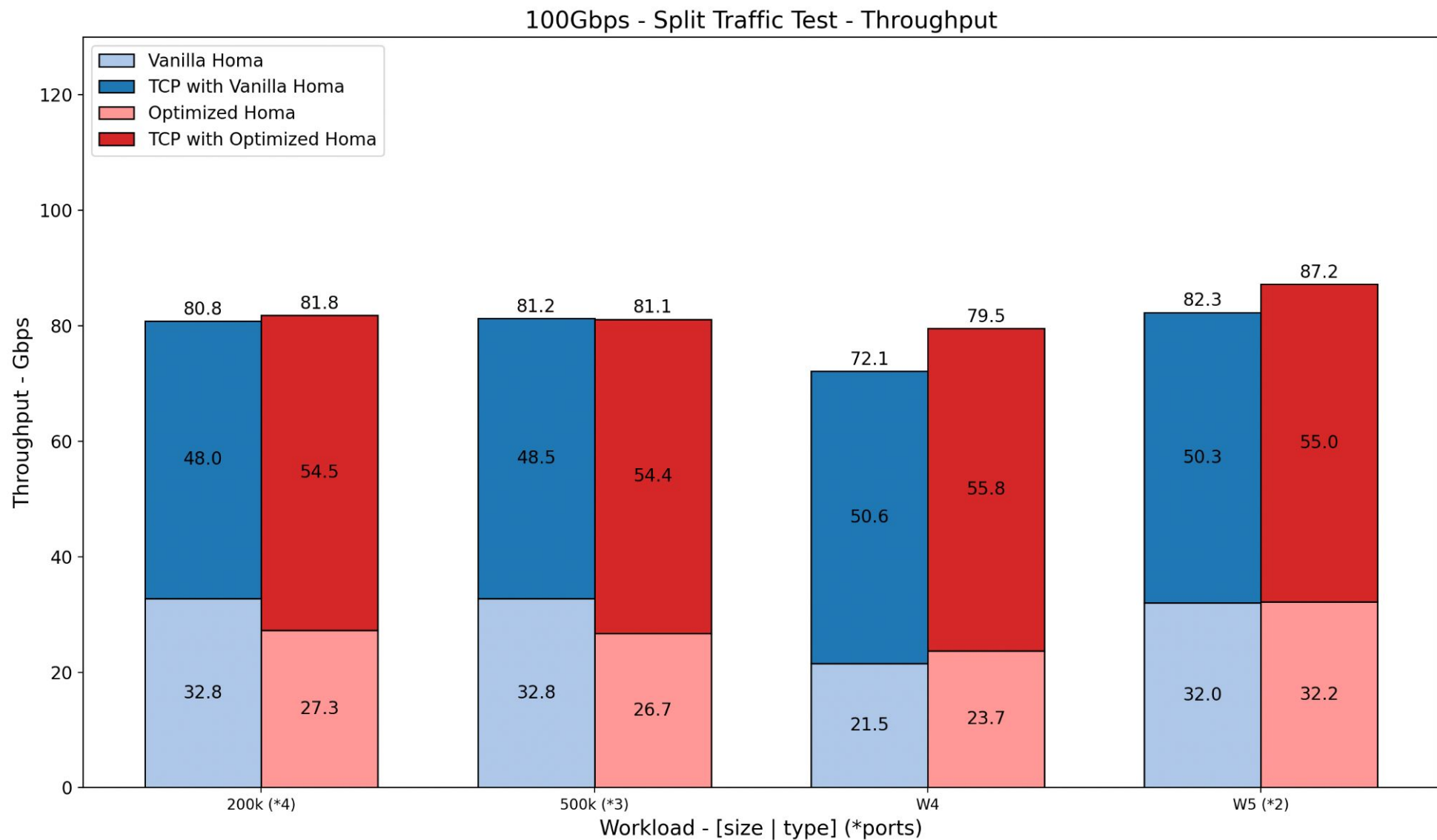
# Long Message Incast - Latency



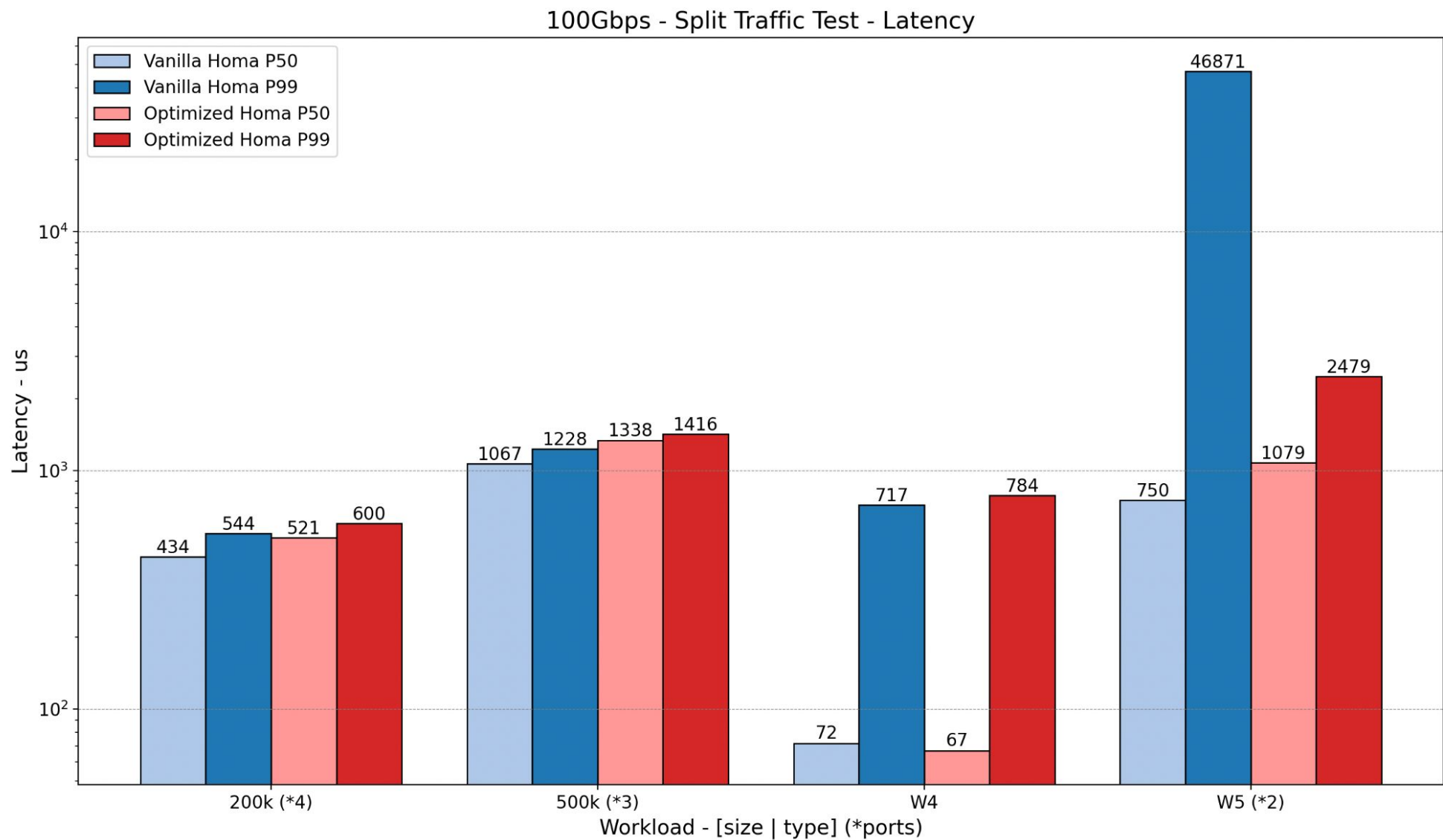
# Split Traffic - Setup



# Split Traffic - Throughput



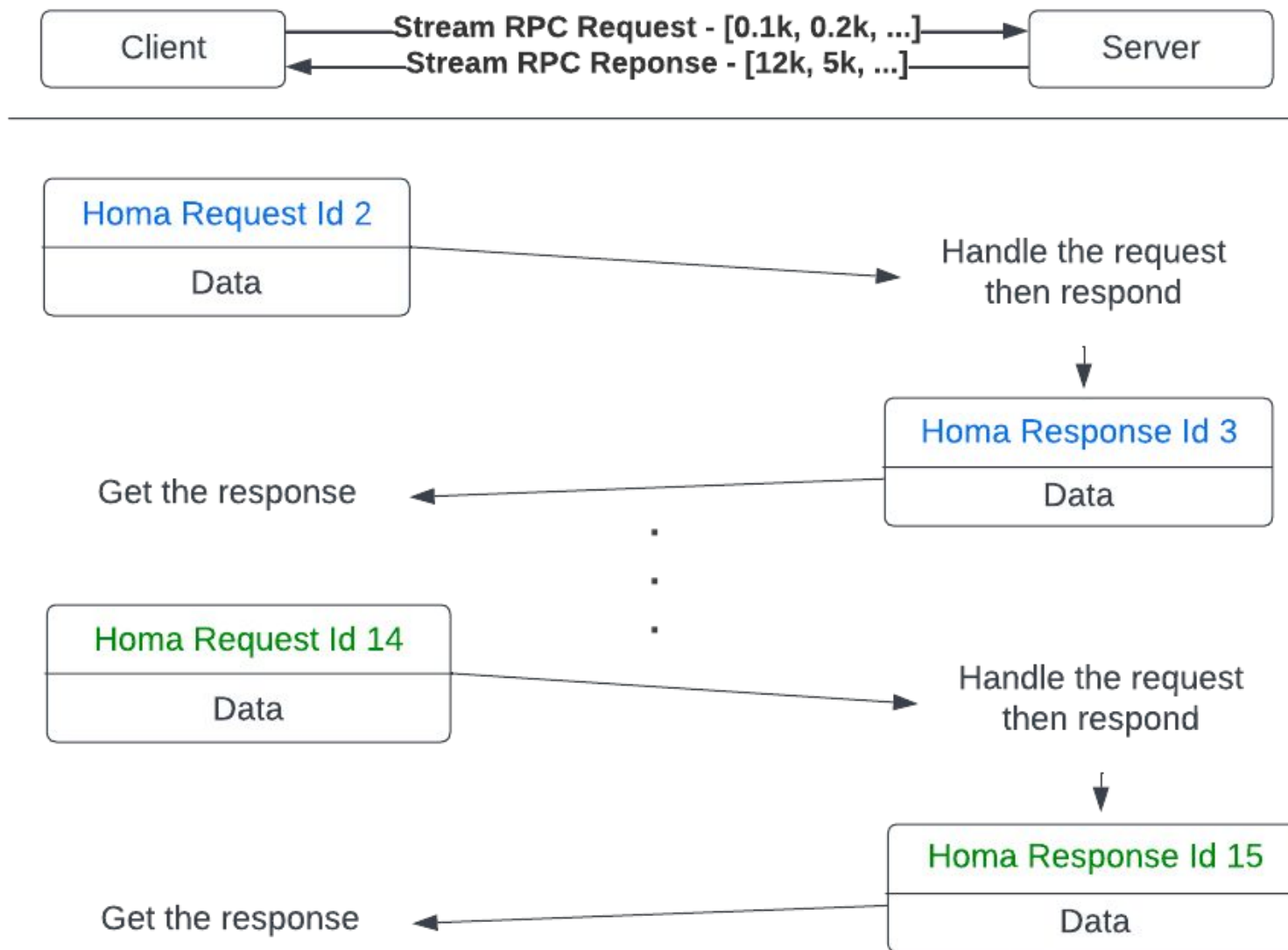
# Split Traffic - Latency



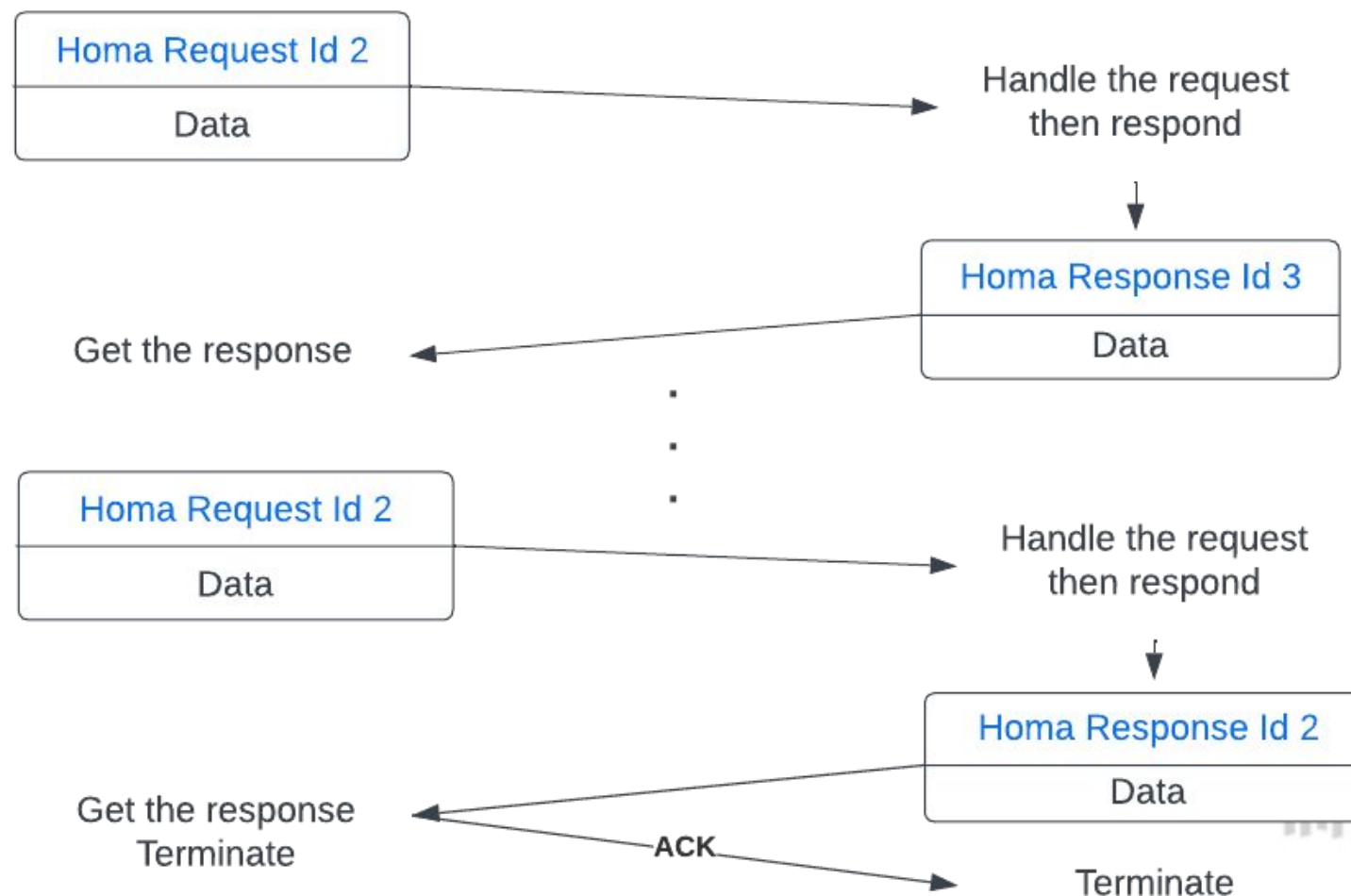


# Homa RPC Streaming Enhancements

# Homa RPC Streaming



# Homa RPC Streaming Enhancements





# Future Improvements



# Future Improvements

- More accurate RTT measurement (Fabric + NIC + software delay)
- Optimize the dynamic window algorithm (Any thoughts?)
- Optimize pacer
- Zero-copy
- More tests for stream RPC enhancements



# Conclusion

1. Homa is a very promising protocol in RPC context
2. Automatic and Dynamic RTTBytes is a better choice
3. Dynamic per peer adjustable window
  - Improve performance on throughput and latency
  - Buffer Overlimit Resilience in Incast
  - Compatibility with TCP Traffic

# Q & A



# References

1. John Ousterhout Stanford University, “A Linux Kernel Implementation of the Homa Transport Protocol”, 2021 USENIX Annual Technical Conference.
2. Radhika Mittal\* (UC Berkeley), Vinh The Lam, Nandita Dukkupati, Emily Blem, Hassan Wassel, Monia Ghobadi\* (Microsoft), Amin Vahdat, Yaogong Wang, David Wetherall, David Zats, “TIMELY: RTT-based Congestion Control for the Datacenter” SIGCOMM ’15 August 17-21, 2015, London, United Kingdom
3. Behnam Montazeri, Yilong Li, Mohammad Alizadeh† , and John Ousterhout Stanford University, +MIT, “Homa: A Receiver-Driven Low-Latency Transport Protocol Using Network Priorities”, SIGCOMM ’18, August 20-25, 2018, Budapest, Hungary
4. John Ousterhout <https://homa-transport.atlassian.net/wiki/spaces/HOMA/pages/262178/Homa+Projects>