Zero Copy Rx with io_uring

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01 Problem statement

02 io_uring primer

03 Design

04 Preliminary results

05 Status + future work

06 Questions + discussions

Agenda

01 Problem statement

01 PROBLEM STATEMENT

DRAM

Linux networking Rx overheads

- Memory and PCIe bandwidth bottlenecks
- Memcpy CPU overheads

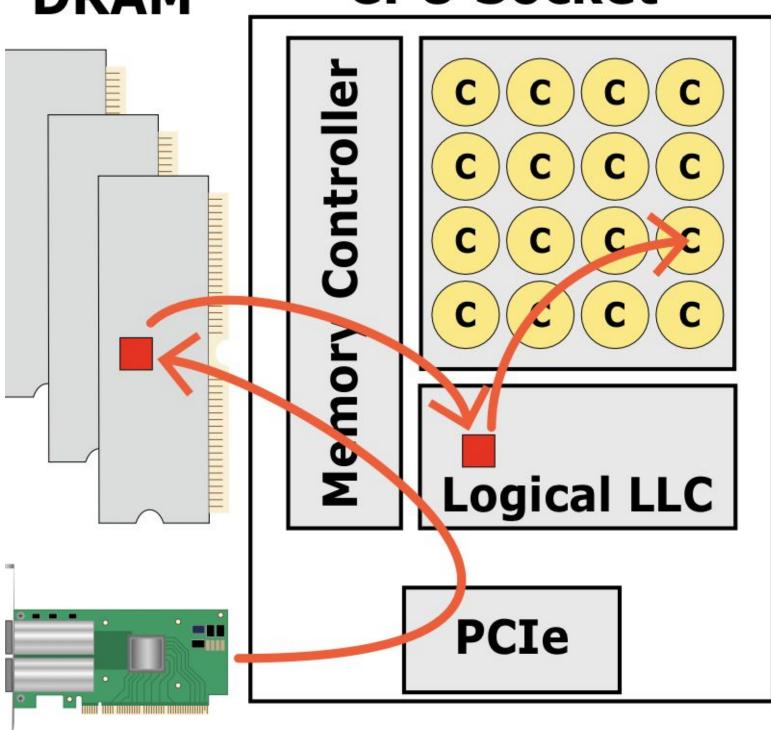


Image credit: A Farshin, A Roozbeh, G Q Maguire Jr., D Kostić. Reexamining Direct Cache Access to Optimize I/O Intensive Applications for Multi-hundred-gigabit Networks.

CPU Socket

Kernel bypass

- High throughput! Low latency!
- **But** libraries and applications *expect* kernel networking stack
- Re-architecting an entire system around kernel bypass is expensive



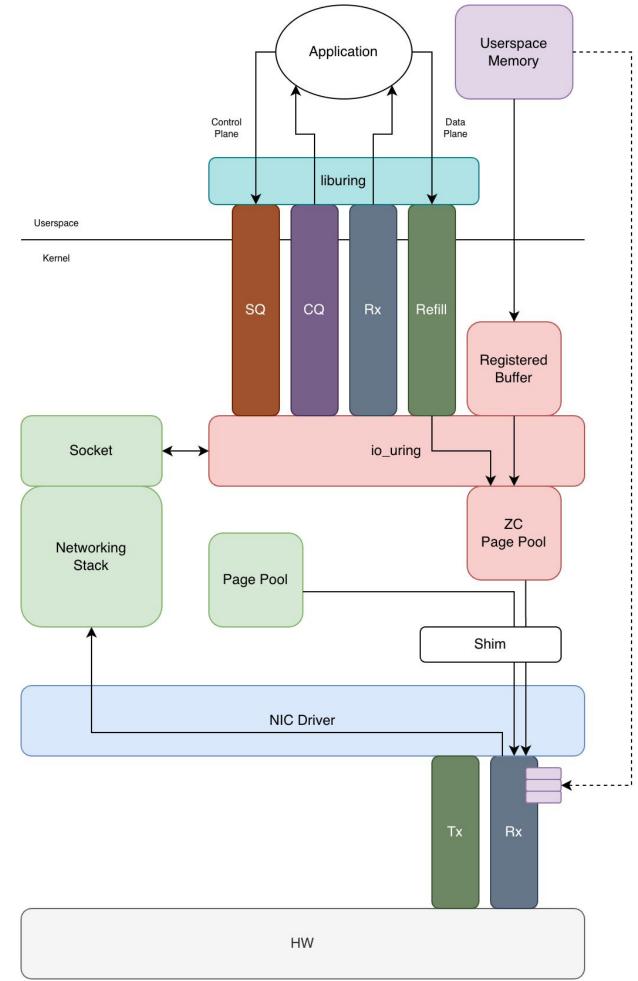


01 PROBLEM STATEMENT

Proposal

Hybrid solution

- Standard control plane using kernel networking stack
- Fast ZC Rx data plane using io_uring
- Two parts:
 - sk_buffs with page frags pointing to userspace pages end up in sockets
 - \circ Read from socket using io_uring



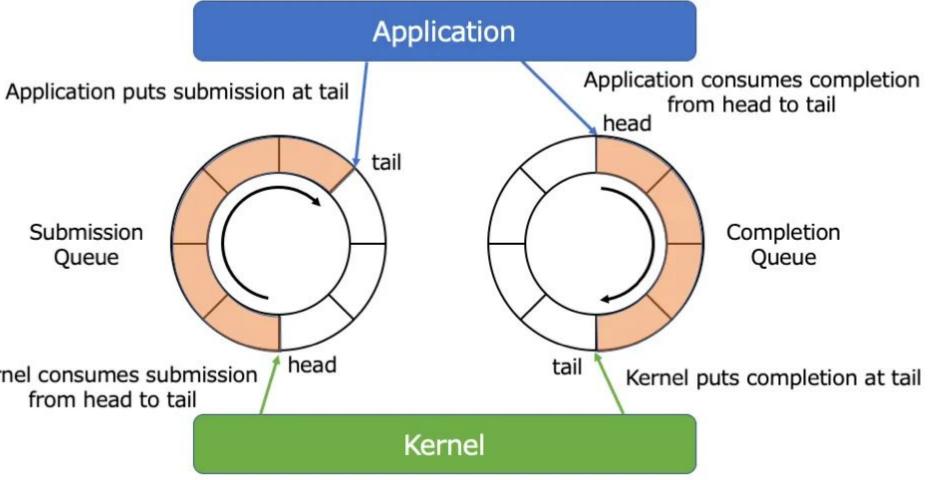
02 io_uring primer

Image credit: https://medium.com/nttlabs/rust-async-with-io-uring-db3fa264 2dd4

io_uring

- Rings shared between kernel and userspace
- Userspace submit requests into Submission Queue (SQ)
- Kernel posts completions into Completion Queue (CQ)
- Kick off work by entering kernel

Kernel consumes submission from head to tail



02 IO_URING PRIMER

02 IO_URING PRIMER

Prepare request

Application puts submission at tail

Submission

Queue

Kernel consumes submission

from head to tail

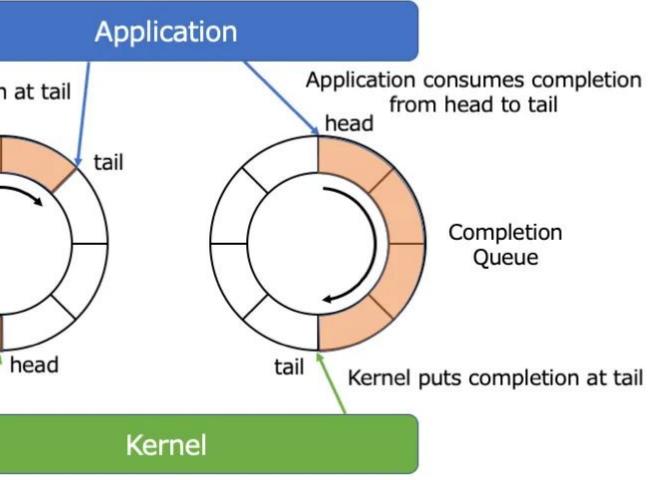
struct io_uring_sqe *sqe;

sqe = io_uring_get_sqe(ring);

io_uring_prep_recv(sqe, sockfd, buf, len, flags);

Note this already moves the SQ tail

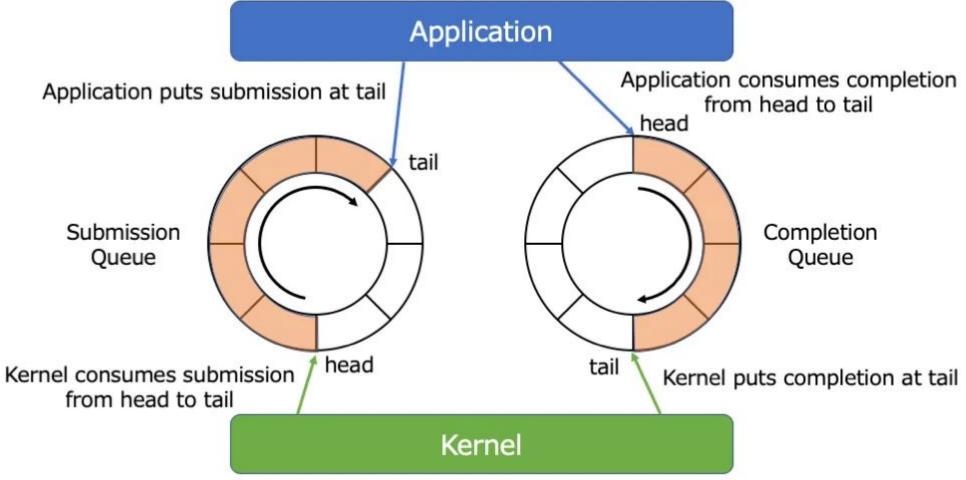
Image credit: https://medium.com/nttlabs/rust-async-with-io-uring-db3fa264 2dd4



Submit

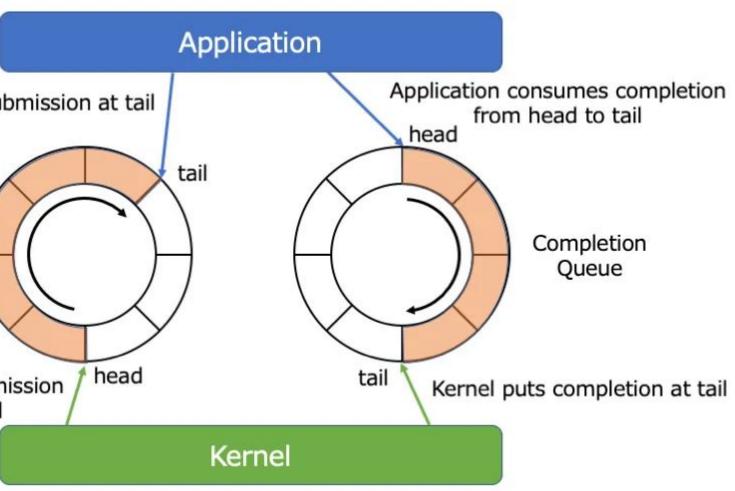
io_uring_submit_and_wait(ring, nr_completions);

Image credit: https://medium.com/nttlabs/rust-async-with-io-uring-db3fa264 2dd4



02 IO_URING PRIMER

Process completions



unsigned head;

int count = 0;

```
io_uring_for_each_cqe(ring, head, cqe) {
```

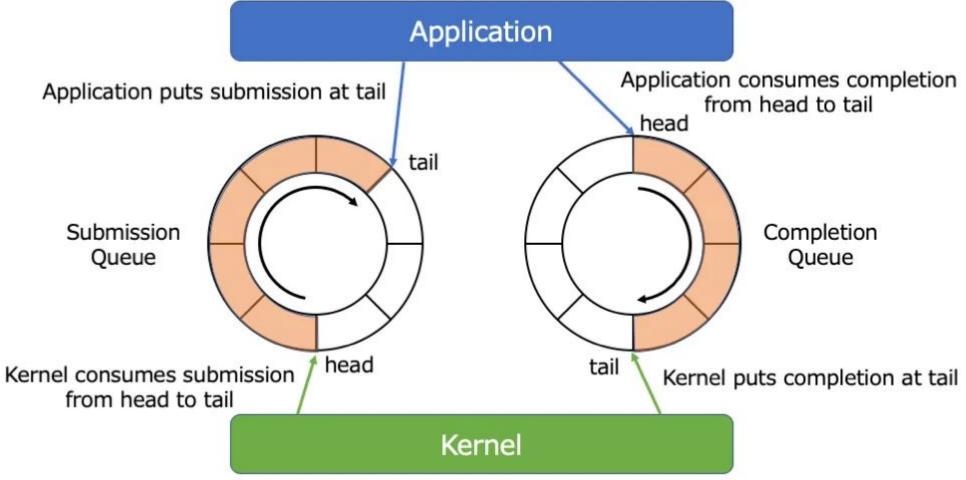
// do stuff

count++;

```
}
```

io_uring_cq_advance(ring, count);

Image credit: https://medium.com/nttlabs/rust-async-with-io-uring-db3fa264 2dd4



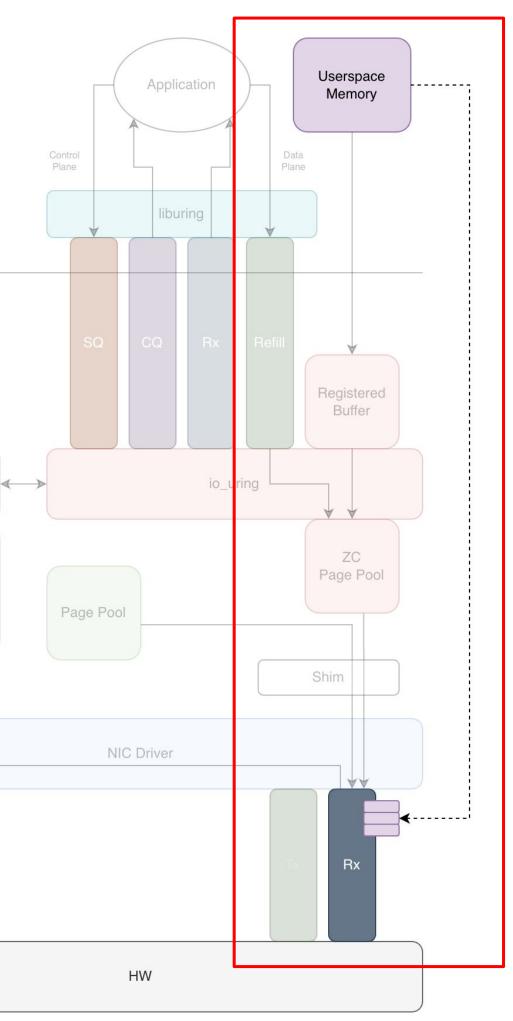
03 Design

Buffer management

- sk_buffs with page frags pointing to userspace pages end up in sockets
- To do this:
 - Fill hardware Rx queue filled with *userspace* pages

Userspace

Kernel

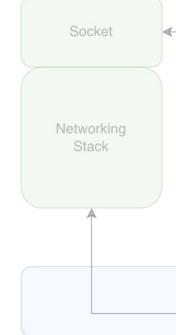


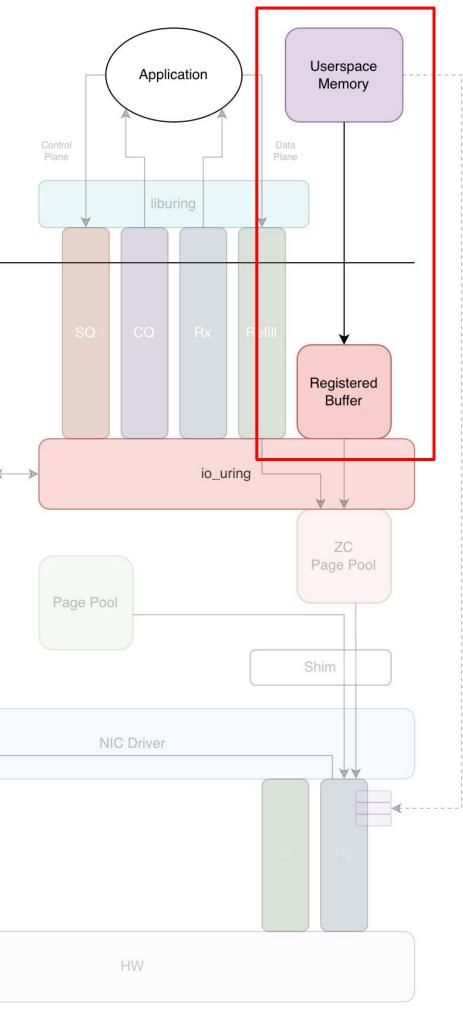
Userspace

Kernel

Buffer management: registration

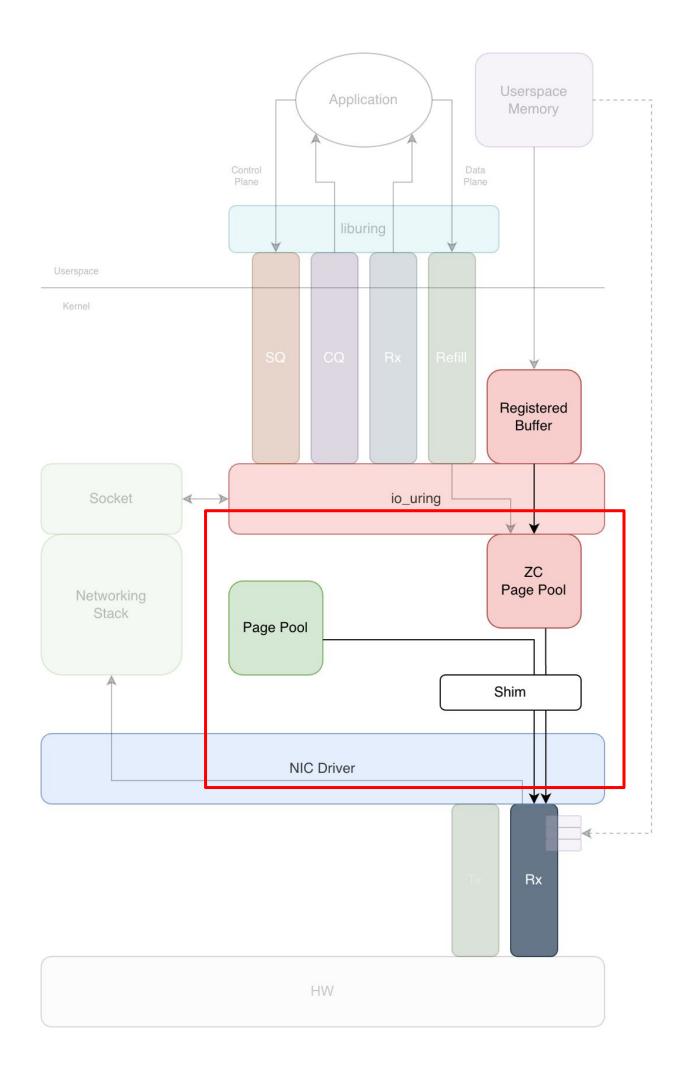
- Register userspace memory with io_uring
- Pin pages
- struct bio_vec bvec[]





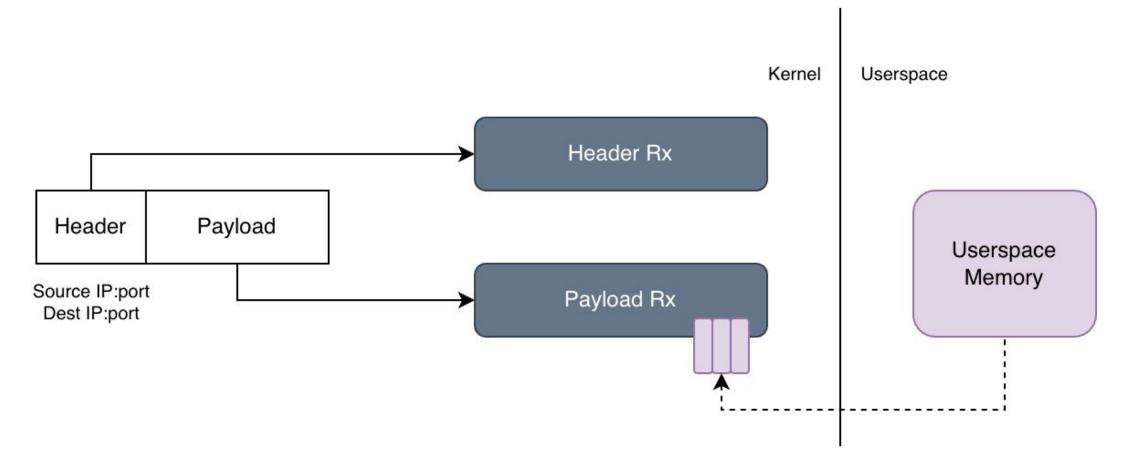
Buffer management: fill HW Rx queue

- Page pool evolving to become generic allocator for NICs
- Add ZC page pool "inspired" by page pool
- Thin shim layer + driver changes



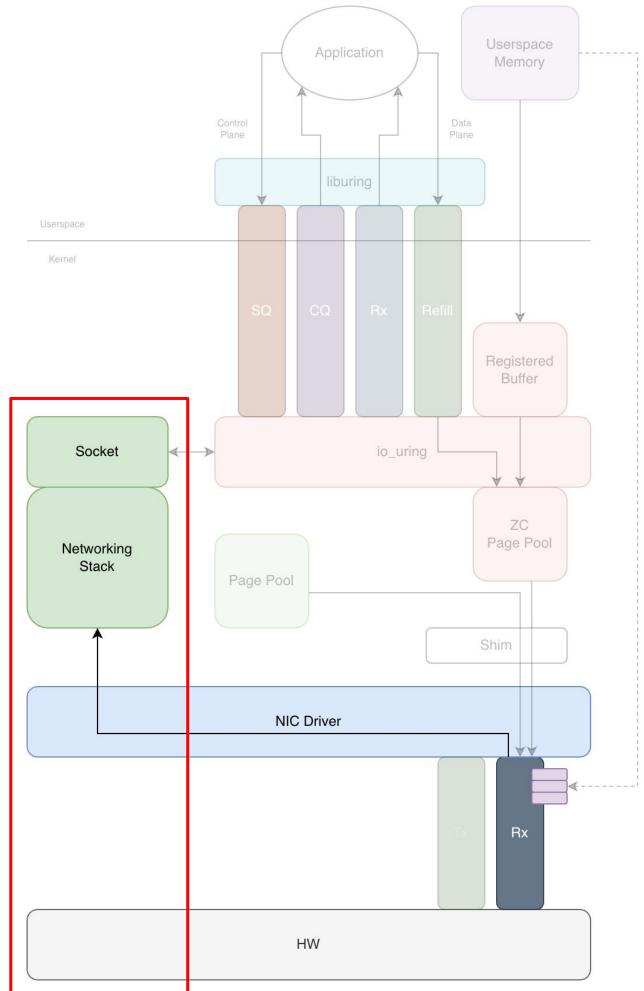
Header splitting + flow steering

- Only want payload
 - Header splitting
- Only want our specific application flows to hit our ZC hardware Rx queues
 - \circ Flow steering
 - RSS



Kernel network stack

- Hardware side fully set up
- Hard IRQs
- Softirq NAPI poll
- Construct sk_buffs
 - Marked as ZC Rx
 - Page frags \rightarrow userspace pages
- Goes through networking stack

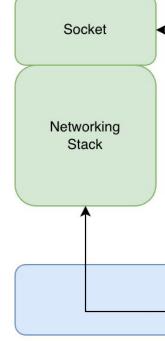


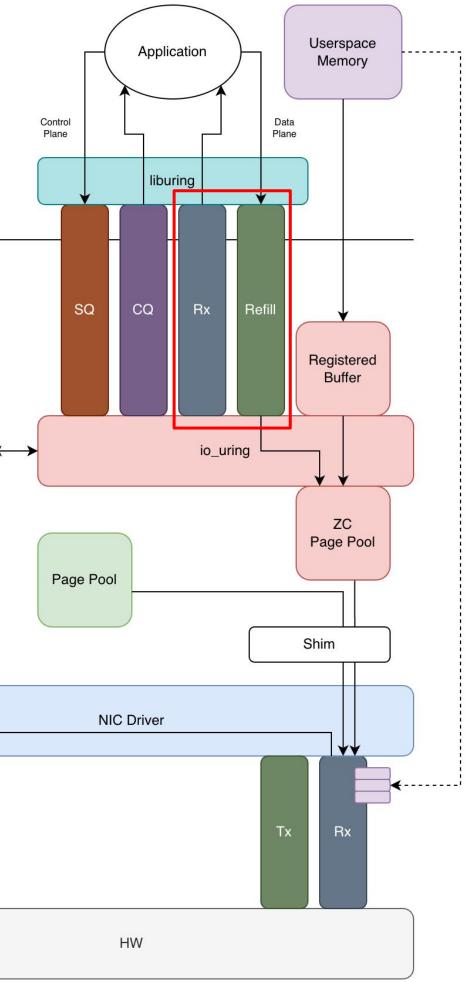
More rings

Userspace

Kernel

- Add two new shared ringbufs to io_uring:
 - Rx queue
 - Refill queue
- One pair per hardware Rx queue



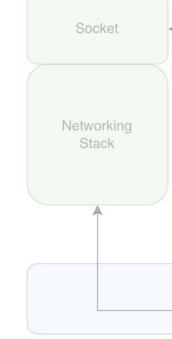


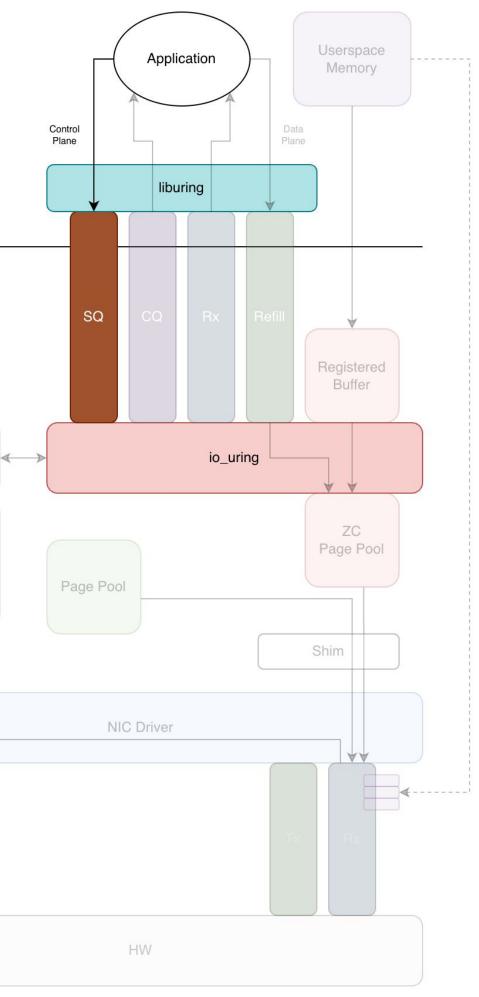
Userspace: submit request

- Submit ZC receive request to io_uring
- Get SQE, prep, and submit

Userspace

Kernel





io_uring: read socket

- Handle ZC receive request
- Read sk_buffs from socket
- No copy payload already in userspace
- Post one ZC Rx queue entry per skb page frag

struct io_uring_rbuf_cqe {

u32 off;

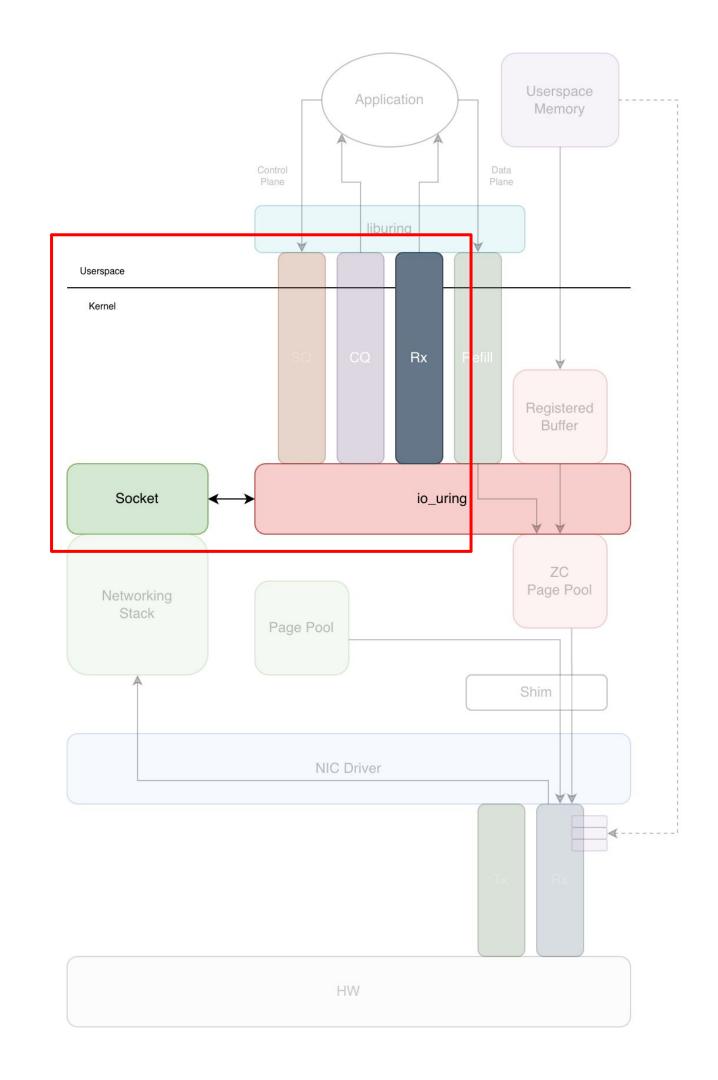
u32 len;

u16 region;

u8 sock;

u8 flags;

}



io_uring: notify userspace

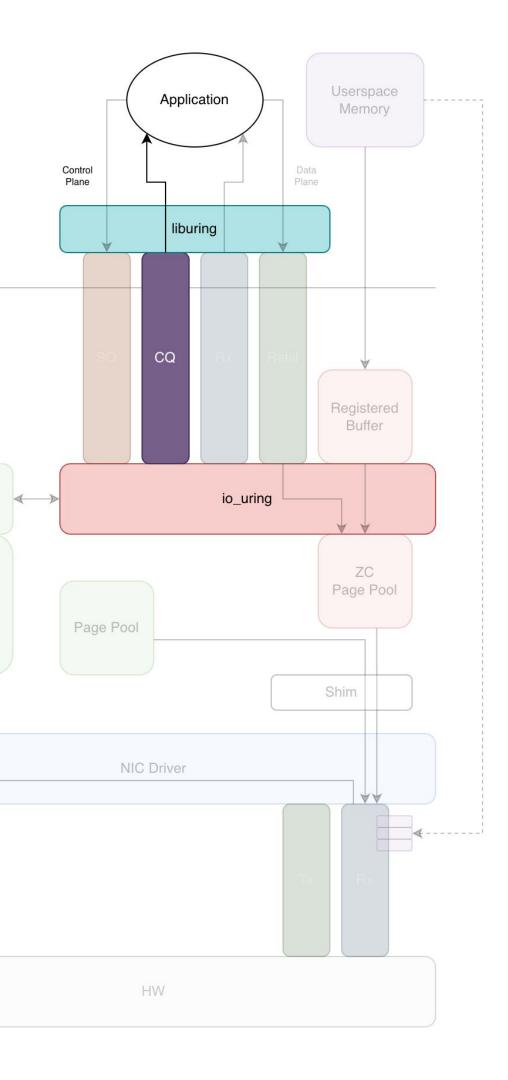
- Post completion event into CQ
- Tells userspace to go look at a ZC Rx queue

Userspace

Kernel

Socket

Networking Stack



Userspace

Kernel

Userspace:

read data

- Look at a ZC Rx queue
- Each entry tells user where the payload is relative to the registered memory region

struct io_uring_rbuf_cqe {

u32 off;

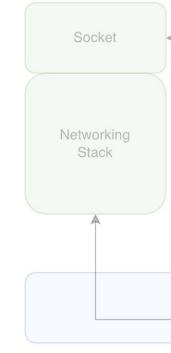
u32 len;

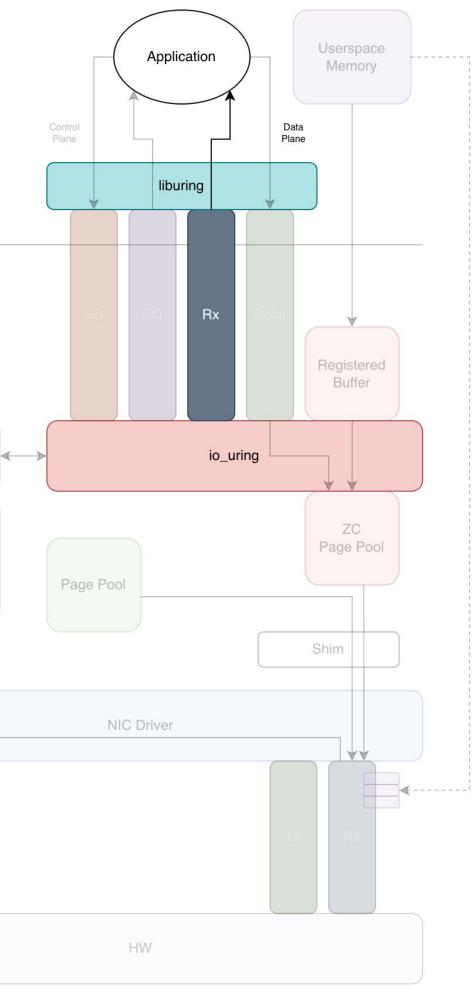
u16 region;

u8 sock;

u8 flags;

}





Userspace: return buffers

- Return buffers to ZC page pool via refill queue
- Eventually used by NIC driver to refill hardware Rx queue

struct io_uring_rbuf_rqe {

u32 off; u32 len;

u16 region;

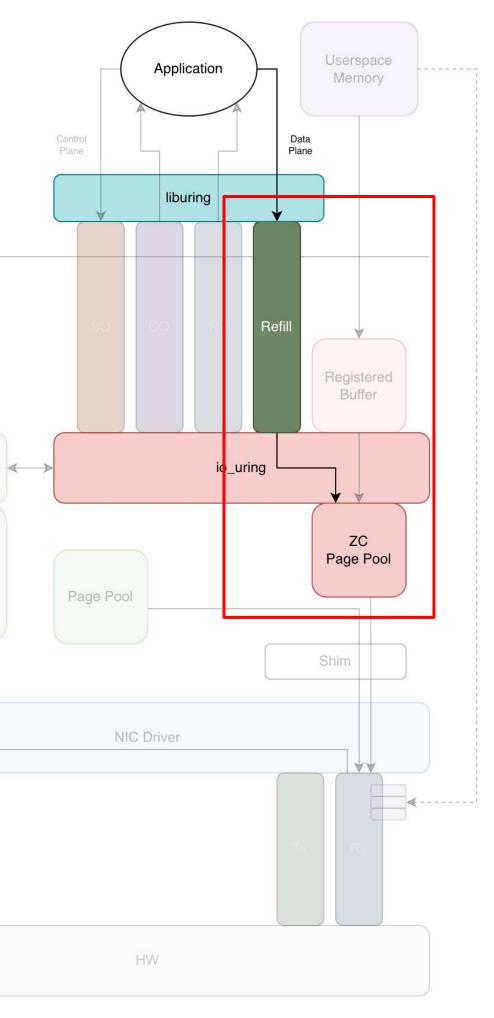
}

Userspace

Kernel

Networking Stack

Socket



04 Preliminary Results

MemBW

Broadcom BCM57504 NIC @ 25 Gbps link

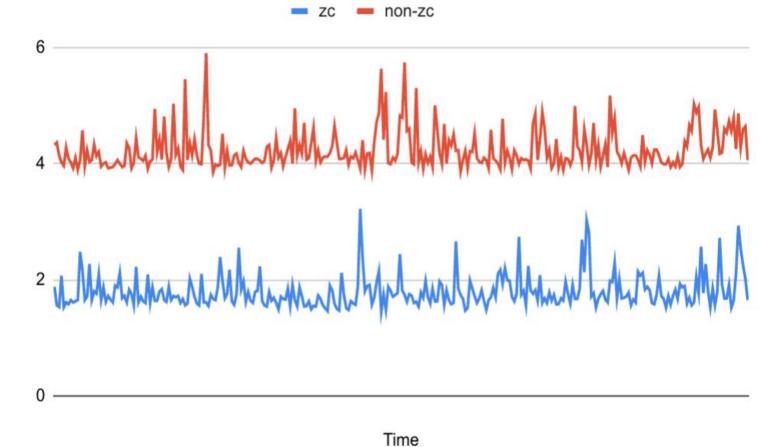
62 GB DRAM

iperf3 + io_uring + ZC Rx

AMD EPYC 7D13

iperf3

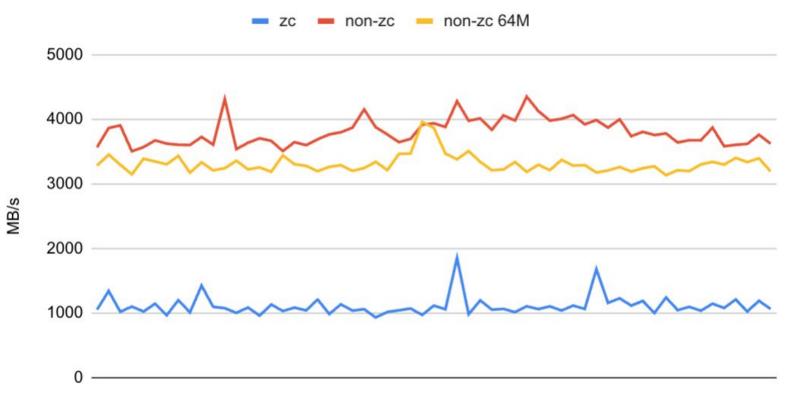
uProf



System Memory Read Bandwidth

GB/s

System Memory Read Bandwidth



MemBW

Broadcom BCM57504 NIC @ 25 Gbps link

62 GB DRAM

iperf3 + io_uring + ZC Rx

Intel Xeon Platinum 8321HC

iperf3

pcm-memory

DDIO is off

Time

05 Status + future work

05 STATUS + FUTURE WORK

Status

- V2 RFC is on the mailing list (netdev + io-uring)
- Hacky veth support if you want to play with the API
- Broadcom bnxt hardware support
- Multi-socket
- Copy fallback

05 STATUS + FUTURE WORK

Future work

- Jakub Kicinski's memory provider API
- Proper test device
 - netdevsim?
- Tying flow steering rules with socket
- Dynamic Rx queue reconfiguration
- Support GPU device memory
 - Using Google's TCP devmem proposal

06 Questions + discussions



07 Appendix

Open questions

- Containers + VMs support?
- TLS + kTLS?
- HugePages?

Copy fallback

- What if we run out of userspace memory allocated for ZC Rx?
- Fill HW Rx queue with kernel pages as before
- When io_uring ZC receive finds sk_buffs with page frags that are not ZC pages, copy into a page from refill queue

Handling errors

- How much to allocate ahead of time?
- What if it runs out?
- What if header splitting fails?
 - Split too little header malformed
 - Split too much payload included
- What if flow steering fails?
 - ZC Rx packet ends up in non-ZC Rx queue
 - Non-ZC Rx packet ends up in ZC Rx queue

Integrating ZC Rx well

- NIC \rightarrow userspace memory is only one hop in a long end to end pipeline
- What if data needs to be modified after ZC Rx? Another copy...
- API need to expose fine control over the placement of data to satisfy constraints e.g. alignment
 - Hardware also needs to support this too