Virtuozzo



Restoring process trees with child-sub-reapers, nested pid-namespaces and inherit-only resources.

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Agenda

- Brief overview of Process tree checkpoint/restore
- Algorithms for simple cases
- Complex case (PR_SET_CHILD_SUBREAPER / CLONE_PARENT)
- Way to solve it: CABA (Closest Alive Born Ancestor)

Brief overview of Process tree checkpoint/restore

- CRIU knows nothing about the history of creation
- Process tree can change when processes fork/exit
- On father exit, children are reparented to reaper:
 - To pid namespace/host init (CLONE_NEWPID)
 - To child-subreaper (PR_SET_CHILD_SUBREAPER)
 - To alive thread of father
- Session (sid) is an example of special resource which:
 - can only be created by specific process session leader (pid == sid)
 - non session leader can only inherit session
 - can inherit "born session" from session leaders

Algorithms for simple cases



Algorithms for simple cases: fork (clone) and setsid

- 1. Traverse process tree and add all non-leaders to the search list.
- 2. While search list is not empty:
 - a. Pop first item
 - b. For session leaders we check their born_sid for others sid
 - c. If checked sid is equal to either parent sid or born_sid continue
 - d. Else set parent born_sid to item's sid
- 3. Traverse process tree forking processes in process tree order
 - a. First fork children with sid or born_sid equal to parent born_sid
 - b. Setsid in parent
 - c. Then fork other children (with sid or born_sid equal to parent sid)

note: this does not handle CLONE_PARENT from session leader

Algorithms for simple cases: fork and setsid

Born sid example:





- Final tree (after reparenting)
- Previous algorithm can't handle born session "collision" here
- Assume no "external" sids
- 🗙 exit



Algorithms for simple cases: + exit and one pidns

- 1. If we have session without leader
 - a. add leader helper as a child of init
- 2. Set born_sid-s in a same manner as in previous algorithm except for init
- 3. For each process which is a child of init
 - a. For session leaders we check their born_sid for others sid
 - b. If checked sid is equal to init sid continue
 - c. Else we search the session leader for the checked sid
 - d. Add a helper process in session leader children
 - e. Move the checked process into helper children
- 4. Fork processes the same way as in previous algorithm

5. At the end we exit from helper processes and get right tree via reparenting note: except for highlighted this is implemented in mainstream CRIU [1], also in CRIU we create helpers for missing processes, e.g. for process groups and accessed /proc/<pid> files

Algorithms for simple cases: + exit and one pidns



note: The tree with helpers can be different to "historical" tree note: If there would be any other "only inheritable" resource which should be inherited from 3 to 6 or from 7 to 9 we would have problem here



note: All pids/sids here and after are shown from root pid namespace



- 1. If we have session without leader
 - a. find proper pid namespace for a leader going up for each zero in multilevel sid
 - b. add leader helper as a child of init of this proper pid namespace
- 2. For each pid namespace init current init
 - a. walk over processes subtree including other inits but excluding their children and setup "born sid" for them (excluding current init)
 - b. when reparenting branches which can't inherit sid from their init create a helper in current pid namespace connected to the session leader
- 3. When forking the tree we create pid namespace inits with CLONE_NEWPID
- 4. When parent pidns is different to child pidns and child is not init we do setns before forking
- 5. After all is forked all helpers exit and we get desired tree



Algorithms for simple cases: + nested pidns (other example)



This one is almost complex

Algorithms for simple cases: + nested pidns (other example)



Algorithms for simple cases: + nested pidns (other example)



- note: we set pids via ns_last_pid or clone3 settid in each pidns level that's why all sids are also right in all pidns levels
- note: in case of nested pid namespaces together with nested user namespaces clone3 settid does not work as it can't set tid on pidns level owned by other userns, I tried to fix it, but got no reaction [2]
- note: process groups become also only inheritable if we enter nested pid namespace
- note: no "external" sids, always do setsid when you enter a container

Algorithms for simple cases: + simple child-subreapers

- When process becomes a child-subreaper (PR_SET_CHILD_SUBREAPER) it can catch reparenting processes similar to init in it's subtree
- "Simple case" is when process becomes child-subreaper just after fork and never disables this (similar to what docker shim does)
- Handle child-subreapers similar to pid namespace inits
 - though no setns analogy available, so we should check born sids for them
 - accurately handle different cases where child-subreaper does not allow altering the tree by moving init or other subreaper branch to it's session leader (session leaders should be created as children of child-subreapers sometimes)
- For more details please just look into Virtuozzo CRIU code [3]
- When I was preparing those pictures I found two issues in algorithm =)

Algorithms for simple cases: + simple child-subreapers



Child-subreaper

Algorithms for simple cases: + simple child-subreapers







Complex case: +child-subreapers (+ CLONE_PARENT)

1. 2 is child-subreaper, 6 exits



2. 2 is not child-subreaper anymore, 3 exits



Complex case: +child-subreapers (+CLONE_PARENT)

Why is it complex?

- Uncertainty which session was created first? sid 2 or sid 4?
- Other similar example from 10 years ago [4]
- Imagine if we have other resource like sid inherited from 2 to 4, it means that we also need to consider this dependency when restoring the process tree
- With child-subreapers and killing the parent we can reparent any process to any of it's ancestors in the same pid namespace as the parent
- CLONE_PARENT is equivalent to child-subreaper reparenting
- Before fixes [5] and [6] child-subreapers were even more complex (non uniform reaper, cross ns reaper)



Complex case: +child-subreapers (+CLONE_PARENT)



note: This can be created similar to previous example

CABA Closest Alive Born Ancestor





CABA: Implementation

- CABA Closest Alive Born Ancestor
- CABDs Closest Alive Born Descendant-s
- Implementation [7] is similar to real_parent/children/sibling (caba/cabds/cabd respectively), except:
 - When process is forked its CABA is set to real_parent (or current for CLONE_PARENT) and process is added to CABDs list of its CABA
 - When father is released (waited/unhashed) its CABDs are moved to father's CABA
 - Except when thread leader of father is available, then thread leader will be new CABA of father CABDs
- "NScaba: ..." line is printed to /proc/<pid>/status

Note: I have an Idea that we can enable CABA only per pidns (recursively) Note: in v2 CABA there was real_parent only (changed in v3) Note: in v2 CABA exited father was handled (changed in v3) CABA



[7] patch to mainstream has selftest with this example

CABA: Restore based on it

- 1. Take CABA process tree
- 2. For each session without session leader add missing session leaders
 - a. find the smallest subtree containing all processes of the session
 - b. from root process in the subtree go up the tree until we are in enough low pid namespace to create the session leader as our child (skip sid zeroes)
 - c. add session leader helper to tree as a child of the found process
 - i. move the found process children which have the session in subtree to the newly created session leader (via helper)
- **3**. Add helper parent processes for each process which parent != caba
- 4. Set "born_sid" to all processes
- Restore processes according to updated CABA tree order handling "born_sid"s same as in first algorithm
- 6. Kill helpers temporary enabling child-subreaper flag for parent in process tree Note: not yet implemented

CABA: Notes

- CABA is self restoring by design
 - if we follow caba tree on restore in CRIU we get the same caba tree
- There are attempts to make a clear mathematical models of processes/resources restore like [8] where order of restore operations is determined based on combined information about all the dependencies between all the resources, unlike CRIU where order is pretty much hardcoded
- But in Linux we generally don't have information about all resources dependencies
 - new resources can appear
 - some dependencies may be unclear or hard to describe in model
- CABA gives an ability to have as much info about process order as we can
- CRIU is not considering per-thread children as on reparent they can move unpredictably
- CABA is not affected by pid-reuse as "historical" tree would be

Links

[1] Mainstream CRIU current implementation:

https://github.com/checkpoint-restore/criu/blob/6128eb6185b4/criu/pstree.c#L681

[2] [PATCH] clone3: add option to change owner of newly created namespaces https://lore.kernel.org/lkml/20210402155131.119872-1-ptikhomirov@virtuozzo.com

[3] Virtuozzo CRIU simple child-subreapers + nested-pidns case:

https://github.com/OpenVZ/vzcriu/blob/1f87bcd638b5/criu/pstree.c#L1573

[4] "Real life task: How to restore process tree in Linux", 2013, Andrei Vagin, in Russian: <u>https://habr.com/en/post/195330</u>

[5] commit "prctl: propagate has_child_subreaper flag to every descendant" https://github.com/torvalds/linux/commit/749860ce2427

[6] commit "exit: fix the setns() && PR_SET_CHILD_SUBREAPER interaction"

https://github.com/torvalds/linux/commit/c6c70f4455d1

[7] [PATCH v3 0/2] Introduce CABA helper process tree

https://lore.kernel.org/lkml/20220908140313.313020-1-ptikhomirov@virtuozzo.com

[8] Efanov, Nikolay. "On Synthetic Process Trees Reconstruction Using Graph-Based Operation Restore Model." 2020 International Conference Engineering and Telecommunication (En&T) (2020): 1-5.