

## DWARF Extensions for Optimized SIMT/SIMD (GPU) Debugging GNU Tools Cauldron @ Linux Plumbers Conference 2021

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## AMD Open Source Debugging Tools Project

- AMD ROCm ROCgdb debugger
  - <u>https://github.com/ROCm-Developer-Tools/ROCgdb</u>
- DWARF Extensions For Heterogeneous Debugging
  - <u>https://llvm.org/docs/AMDGPUDwarfExtensionsForHeterogeneousDebugging.html</u>
- User Guide for AMDGPU Backend: DWARF Debug Information
  - https://llvm.org/docs/AMDGPUUsage.html#dwarf-debug-information

## Heterogeneous Computing Debugging Challenges

- GPUs and other heterogeneous computing devices have:
  - Multiple memory address spaces
  - Many wide vector registers
  - Many scalar registers
  - Language threads of execution => lanes of a SIMD/SIMT execution model



### Heterogeneous Computing Debugging Challenges (cont.)

- Variables more often spread across pieces of different storage kinds
- SIMD/SIMT execution needs runtime selection of pieces of vector registers
- Complex expressions benefit from factorization of location definitions
  - Currently not possible => duplicate parts of location definitions

### **Existing DWARF 5 Limitations**

- Unable to describe variables in combinations of parts of registers => no static or runtime indexing
- Some features only work when located in memory => type attribute expressions requiring a base object
- DWARF procedures can only accept global memory address arguments
- No vector base types needed to describe vector registers
- Cannot create memory locations in address spaces
- CFI does not allow composite locations
- CFI does not support address spaces
- Bit field offsets are not supported for all location kinds

## How to fix this?

- Explored numerous approaches to overcome limitations
  - Approach chosen was simplest and provided most benefits
- Based on:
  - Generalizing execution model
  - Composable and consistent operations
- Results in small number of new operations that compose generally
- As opposed to adding many specialized operations and rules
  - Causes contextual semantics and corner cases
  - Harder for compilers and debuggers
- Major aspect is to allow locations to be manipulated on the stack

## **Main Goals**

- Upstreamable:
  - GDB debugger
  - LLVM compiler
  - GCC compiler
- Supportable by other tools:
  - TotalView debugger
- Backwards compatible with DWARF 5
- Support optimized code for GPUs
- Benefit non-GPU targets too

## What Is DWARF?

- A standard way to specify debug information
  - Describes source language entities: compilation units, functions, types, variables, etc.
  - Embedded in sections of code object executables
- Maps source program language entities to the hardware representation:
  - Program counter <=> source line
  - Source function => entry point program counter
  - Source language variable => location at a particular program counter
  - Source function call stack virtual unwinding => locations of callee saved registers
  - Etc.

### **DWARF Expressions**

- Great diversity in locations of entities
- Locations involve runtime values
- Variable location could be:
  - In register
  - At memory address
  - At Stack Pointer + offset
  - Optimized away known value => such as compile time value
  - Optimized away unknown value => such as unused variables
  - Spread across combination of locations
  - At memory address, but also transiently loaded into registers
- Consequently, DWARF includes rich expression language
  - Expression evaluated on simple stack machine
  - Expression evaluated in a context => value/location result kind, process/thread/lane/frame/pc, initial stack, etc.
  - Some context defined by place expression used => result kind, initial stack, etc.

#### **DWARF 5: Dynamic Array Size (1 of 3)**

Expression DWARF PC DW\_OP\_regval\_type SGPR0 Generic DW\_OP\_deref

#### **Example: Runtime size of a dynamic array**

Context result kind => value

10

• Expression evaluated one operation at a time using a stack



Stack

Result:Value

Context

#### **DWARF 5: Dynamic Array Size (2 of 3)**



- DW\_OP\_regvalue\_type => reads register and pushes value on stack
- Each stack element is a value and associated type
- Type must be a DWARF base type => specifies encoding, byte ordering, and size of value
  - Defaults to Generic type => architecture specific integral encoding/byte ordering, the size of global memory address



#### **DWARF 5: Dynamic Array Size (3 of 3)**



• Value result kind => result is the top stack value



### **DWARF 5: Variable Location in Register (1 of 2)**

Expression

Stack

Location

Context

DWARF PC----

DW\_OP\_regx SGPR0

Result:Location

#### **Example: Source variable located in a register**

- Context result kind => location
  - Note: DWARF uses term location description

Source Program	
Variable ∭	
Hardware	$\overline{\mathbb{Q}}$
SGPR 0	

### **DWARF 5: Variable Location in Register (2 of 2)**



- DW\_OP\_regx => creates register location
- DWARF conceptually has a separate location area
  - Does not use the stack
- Location result kind => result is the location area

Source Program	
Variable	
Hardware	$\overline{\mathbf{V}}$
SGPR 0	

#### **DWARF 5: Variable Location in Memory (1 of 4)**

Expression DWARF PC DW\_OP\_regval\_type SGPR0 Generic DW\_OP\_plus\_uconst 0x10 Stack

Location

Context

Result:Location

Example: Source variable in stack frame memory at address stack pointer + 0x10



### **DWARF 5: Variable Location Memory (2 of 4)**



 DW\_OP\_regvalue\_type => pushes value read from stack pointer register
 Source Progra



### **DWARF 5: Variable Location Memory (3 of 4)**



- Variable location is 0x10 bytes from the base of the stack frame
- DW OP plus uconst => pop value, add 0x10, and push result



### **DWARF 5: Variable Location Memory (4 of 4)**



- Location area empty when location result kind => convert top stack element to memory location
  - Use value as global memory address



### **DWARF 5: Variable Spread Across Different Locations (1 of 7)**



# Example: Source variable that is partly in a register, partly undefined, and partly in memory

- Composite location => 1 or more parts
  - Each part specifies a location and number of bytes used from it



#### **DWARF 5: Variable Spread Across Different Locations (2 of 7)**



DW\_OP\_regx => creates register location



#### **DWARF 5: Variable Spread Across Different Locations (3 of 7)**



- DW OP piece => first use creates incomplete composite location
  - Location in location area used in first part
  - Size 4 indicates number of bytes used from beginning of part's location



### **DWARF 5: Variable Spread Across Different Locations (4 of 7)**



- DW\_OP\_piece => subsequent use adds part to already created incomplete composite location
  - Parts form a contiguous set of bytes
  - If no other location in location area, and no value on stack => part implicitly the undefined location
  - 2 indicates there are 2 undefined bytes
- Undefined location => used to indicate part that has been optimized away



#### **DWARF 5: Variable Spread Across Different Locations (5 of 7)**



 DW\_OP\_bregx => read register as Generic type, add 0x10, and push value



#### **DWARF 5: Variable Spread Across Different Locations (6 of 7)**

Stack



- DW\_OP\_piece => add part to already created incomplete composite location
  - If no other location in location area, but value on stack => part is memory location with address popped from stack
  - 2 indicates there are 2 bytes from memory





#### **DWARF 5: Variable Spread Across Different Locations (7 of 7)**



Incomplete composite location implicitly converted to complete composite location



2

Address 0x0a3c0f10

### **DWARF 5: Offsetting a Composite Location (1 of 2)**



#### Example: Offsetting a composite location not supported

- Extend previous example to offset location built so far
- Variable Location in Memory example used DW\_OP\_plus => convenient way to offset memory address

## DWARF 5: Offsetting a Composite Location (2 of 2)



- However, DW OP plus cannot be used to offset a composite location => it only operates on the stack
- Compiler would need to make a different composite location => starting at the part corresponding to offset

```
DW_OP_piece 1
DW_OP_bregx SGPR0 0x10
DW_OP_piece 2
```

• Operations on values are not composable with locations

## What Is A Location?

- Location storage is contiguous linear organization of certain number of bytes
- All location kinds can be viewed the same way:
  - Global memory => linear stream of bytes of the architecture's address size
  - Register => linear stream of bytes of the size of each architecture's register
  - Composite location => linear stream of the bytes defined by the parts
  - Implicit location => linear stream of bytes of the value using the type's byte ordering
  - Undefined location => infinitely sized storage where every byte is undefined
- A location is comprised of:
  - A kind (memory, register, etc.)
  - A reference to a specific location storage of that kind
  - A zero-based offset within the location storage

#### **Stack Location Operations**

- If location could be allowed on the stack:
  - Define new operations to work on locations in compossible manner
  - Example: new DW\_OP\_LLVM\_offset => updates offset of any location kind
- Existing operations can be generalized => act on locations of any kind
  - Example: DW\_OP\_deref => pop a location (rather than memory address value), read it
  - Backwards compatibility via implicit conversions
- Key part of extension is allowing locations on stack
  - DWARF 5 expressions can be evaluated unchanged and yield same results
  - Extension allows greater expressiveness => see following examples

### Extension: Variable spilled to part of VGPR (1 of 3)

Expression

Stack

Context

DWARF PC DW\_OP\_regx VGPR0 DW\_OP\_LLVM\_offset\_uconst 20

Result:Location

# Example: Compiler spills variable stored in SGPR register to fixed lane of VGPR register

- In this case lane 5 of VGPR0 => each VGPR lane is 4 bytes
  - So index is 5 \* 4 = 20

Source F	Progra	am						
Variable	(111)							
Hardwar	e			ł	7			
SGPR 0			<u>spi</u>	lled				
VGPR 0	Lane 0	Lane 1	Lane 2	Lane 3	Lane 4	Lane 5	•••	Lane 63

### Extension: Variable spilled to part of VGPR (2 of 3)

DWARF PC DW\_0P\_regx VGPR0 DW\_0P\_LLVM\_offset\_uconst 20

Register				
Offset	Number			
0	VGPR0			
Тс	op			

Stack

Context

Result:Location

• DW OP regx => now pushes location on stack with byte offset of 0



## Extension: Variable spilled to part of VGPR (3 of 3)



- DW\_OP\_LLVM\_offset => can offset register location
- DWARF 5 does not support specifying register offset => can only have locations starting at beginning
- Defining register names for every part of every register => not practical for GPUs due to sheer number
- Separate register names would not allow computed runtime indexing of register parts
- GPU compilers frequently locate variable in parts of the numerous wide vector registers
  - Especially in optimized code to avoid memory accesses
  - Runtime indices used to support SIMT execution model
- Result is top stack entry



## Extension: Variable across multiple VGPRs (1 of 14)

Stack

Expression DWARF PC DW\_OP\_regx VGPR0 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_regx VGPR1 DW\_OP\_regx VGPR1 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_LLVM\_offset DW\_OP\_piece 4

# Example: Source variable located across a SIMT lane of multiple VGPR registers

- GPU compiler maps language threads to VGPR lanes in SIMT manner
- Thread's variable is spread across the same lane of multiple VGPRs
- Context specifies lane => lane corresponds to focused source thread

Result:Location Lane:5

Context



### Extension: Variable across multiple VGPRs (2 of 14)

#### Expression

DWARF PC DW\_OP\_regx VGPR0 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_piece 4 DW\_OP\_regx VGPR1 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_piece 4

Register				
Offset	Number			
Θ	VGPR0			
T	0p			

Stack

Context

Result:Location



#### Extension: Variable across multiple VGPRs (3 of 14)

#### Expression

DW\_OP\_regx VGPR0 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_piece 4 DW\_OP\_regx VGPR1 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_LLVM\_offset DW\_OP\_piece 4



Stack

Context

Result:Location

Lane:5



DW\_OP\_LLVM\_push\_lane => pushes value of context's lane



### Extension: Variable across multiple VGPRs (4 of 14)

#### Expression

DW\_OP\_regx VGPR0 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_piece 4 DW\_OP\_regx VGPR1 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_mul DW\_OP\_mul DW\_OP\_mul DW\_OP\_mul DW\_OP\_mul

Regi	ster			
Offset	Number			
0	VGPR0			
Value				
Type	Value			
Generic	5			
Va	lue			
Туре	Value			
Generic	4			
T				

Stack

Context

Result:Location

Lane:5

 Each VGPR lane is 4 bytes => lane must be multiplied by 4 to get register byte index

				Focus	ed Thr	ead	
9			Ł	Ļ			
				Focu	sed La	ne	
					Ţ		
Lane 0	Lane 1	Lane 2	Lane 3	Lane 4 L	ane 5		Lane 63
	e Lane 0	Lane 0 Lane 1	Eane 0 Lane 1 Lane 2	Lane 0 Lane 1 Lane 2 Lane 3	Focu	Focused Lane 0 Lane 1 Lane 2 Lane 3 Lane 4 Lane 5	Focused Lane

#### Extension: Variable across multiple VGPRs (5 of 14)

#### Expression

DW\_OP\_regx VGPR0 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_regx VGPR1 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_LLVM\_offset DW\_OP\_piece 4

Register				
Offset	Number			
0	VGPR0			
Value				
Type	Value			
Generic	20			
Тор				

Stack

Context

Result:Location



#### Extension: Variable across multiple VGPRs (6 of 14)

#### Expression

DW\_OP\_regx VGPR0 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_piece 4 DW\_OP\_regx VGPR1 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_LLVM\_offset DW\_OP\_piece 4

Register				
Offset	Number			
20	VGPR0			
Тс	op			

Stack

Context

Result:Location



### Extension: Variable across multiple VGPRs (7 of 14)



a global memory location



- Checks if the top stack element is an incomplete composite location
  - If not, creates new incomplete composite location with no parts
  - Otherwise, pops previously created incomplete composite location
- Adds new part to incomplete composite location and pushes on stack

#### Extension: Variable across multiple VGPRs (8 of 14)

#### Expression

DW\_OP\_regx VGPR0 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_regx VGPR1 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_LLVM\_offset DW\_OP\_piece 4



Stack

Context

Result:Location



#### Extension: Variable across multiple VGPRs (9 of 14)

#### Expression

DW\_OP\_regx VGPR0 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_regx VGPR1 DW\_OP\_regx VGPR1 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_piece 4

Con	posit	e(incomple	ete)
Offset	I	Parts	
0	Size	Loca	tion
		Regi	ster
	4	Offset	Number
		20	VGPR0
	Reg	jister	
Offs	set Number		
0		Ve	PR1
	V	alue	
Тур	be	Va	lue
Gene	ric		5
		Тор	

Stack

Context

Result:Location

Lane:5



DWARF PC

#### Extension: Variable across multiple VGPRs (10 of 14)

#### Expression

DW\_OP\_regx VGPR0 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_regx VGPR1 DW\_OP\_regx VGPR1 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_LLVM\_offset DW\_OP\_piece 4

Con	posit	<b>e</b> (incompl	ete)	
Offset		Parts	;	
0	Size	Loca	ntion	
		Regi	ster	
	4	Offset	Number	
		20	VGPR0	
Register				
Off	set Number			
0	VGPR1		GPR1	
	V	alue		
Тур	oe	Va	alue	
Gene	Generic		5	
	V	alue		
Тур	e e	V.	alue	
Gene	ric		4	
		Тор		

Stack

Context

Result:Location



### Extension: Variable across multiple VGPRs (11 of 14)

#### Expression

DW\_OP\_regx VGPR0 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_regx VGPR1 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_LLVM\_offset DW\_OP\_piece 4



Stack

Context

Result:Location



#### Extension: Variable across multiple VGPRs (12 of 14)

#### Expression

DW\_OP\_regx VGPR0 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_regx VGPR1 DW\_OP\_LLVM\_push\_lane DW\_OP\_uconst 4 DW\_OP\_mul DW\_OP\_mul DW\_OP\_LLVM\_offset DW\_OP\_piece 4



Stack

Context

Result:Location



### Extension: Variable across multiple VGPRs (13 of 14)

#### Expression

Stack

Context

DW\_OP\_regx VGPR0 Composite(incomplete) DW\_OP\_LLVM\_push\_lane Offset Parts DW\_OP\_uconst 4 Location 0 Size DW OP mul Register Offset Number DW OP LLVM offset 4 VGPR0 20 DW OP piece 4 DW\_OP\_regx VGPR1 Register Offset Number DW\_OP\_LLVM\_push\_lane 4 VGPR1 20 DW OP uconst 4 DW\_OP\_mul Тор DW OP LLVM offset DW\_OP\_piece 4

Result:Location

Lane:5

 DW\_OP\_piece => pops location and adds as new part to incomplete composite location on top of stack

Variable					Focu	sed Thre	ad
Hardwar	е			Ł	ļ		
					Foc	used Lan	е
						Ļ	
	l ana 0	Lane 1	Lane 2	Lane 3	Lane 4	Lane 5	Lane 63
	Lane U						
VGPR 0							•••

DWARF PC-

### Extension: Variable across multiple VGPRs (14 of 14)

#### Expression

If top of stack is incomplete composite location =>

implicitly converted to complete composite location

Stack

Context

Result:Lo	catior
-----------	--------

Lane:5







### Extension: Variable across multiple kinds of locations (1 of 7)



Memory

 Same as previous example, except last 4 bytes are from memory and a constant value

. . .

### Extension: Variable across multiple kinds of locations (2 of 7)



Memory

. . .

### Extension: Variable across multiple kinds of locations (3 of 7)



#### Extension: Variable across multiple kinds of locations (4 of 7)

VGPR0



Last 2 bytes are constant value 0xf00d ٠



### Extension: Variable across multiple kinds of locations (5 of 7)



- Pops value
- Creates implicit location storage using value's base type size and byte order •
- Pushes implicit location referencing implicit location storage

. . .

. . .

0xbeef

Memory

•

### Extension: Variable across multiple kinds of locations (6 of 7)



. . .

. . .

0xbeef

Memory

٠ location

### Extension: Variable across multiple kinds of locations (7 of 7)



#### **Extension: Address Spaces (1 of 5)**

Expression DWARF PC DW\_OP\_regval\_type SGPR0 Generic DW\_OP\_uconst 1 DW\_OP\_LLVM\_form\_aspace\_address DW\_OP\_LLVM\_offset 0x10 Stack

Context

Result:Location

# Example: Source variable in stack frame address space memory at address stack pointer + 0x10

- Devices can have multiple hardware supported address spaces
  - Specific hardware instructions to access address spaces
- DWARF 5 DW\_OP\_xderef => dereferences a memory address using an address space
  - No way to create address in a specific address space
  - No way to include address space memory locations in parts of composite locations



#### **Extension: Address Spaces (2 of 5)**



- GPUs use separate address space for per lane managed storage
   => used by stack pointer
- DW\_OP\_regval\_type => push stack pointer address



#### **Extension: Address Spaces (3 of 5)**



Value	
Type	Value
Generic	0x0a3c0f00
Value	
_	
Туре	Value
<i>Type</i> Generic	Value 1

Stack

Context

Result:Location

- DW\_OP\_uconst => push address space number
  - Architecture defines numbers => address space 1 is per lane memory



### **Extension: Address Spaces (4 of 5)**





- DW\_OP\_LLVM\_form\_aspace\_address => pops value and address space number, and pushes memory location which includes the address space
  - Each address space is a separate memory location storage
  - All operations on locations work with memory locations regardless of address space
  - Every architecture defines address space 0 => default global memory address space
- Generalization avoids creating specialized operations to work with address spaces



#### **Extension: Address Spaces (5 of 5)**



Context

Result:Location

- The source variable is at byte 0x10 in the frame
- DW\_OP\_LLVM\_offset => works the same with memory locations that have an address space



#### Extension: Bit Offsets (1 of 4)

Expression

Stack

Context

DWARF PC DW\_OP\_regx SGPR3 DW\_OP\_uconst 20 DW\_OP\_LLVM\_bit\_offset

Result:Location

#### Example: Variable in bit field of a register

- Locations specify an offset within associated location storage => extension allows bit offsets
  - DWARF 5 does not support general bit offset => only supports bit fields in composites with DW\_OP\_bit\_piece
  - DWARF 5 only supports locations that start at the beginning of a register
- Supporting bit offsets benefits all targets

Source Program	
Variable	
Hardware	$\overline{\mathbf{V}}$
SGPR 3 SGPR 3	

#### Extension: Bit Offsets (2 of 4)

Expression

DWARF PC DW\_OP\_regx SGPR3 DW\_OP\_uconst 20 DW\_OP\_LLVM\_bit\_offset

Register	
Offset	Number
0 bits	SGPR3
Тор	

Stack

Context

Result:Location

Source Program	
Variable	
Hardware	Ų
SGPR 3	
20 010	

#### **Extension: Bit Offsets (3 of 4)**



Register	
Offset	Number
0 bits	SGPR3
Value	
Туре	Value
Generic	20
Тор	

Stack

Context

Result:Location

- DW\_OP\_uconst => push bit offset on stack
  - This could also be a runtime calculation

Source Program	
Variable	
Hardware	Ţ
SGPR 3	

### Extension: Bit Offsets (4 of 4)





- DW\_OP\_LLVM\_bit\_offset => pop value and location, update location's offset using value as a bit offset, push updated location
- Bit ordering, like byte ordering, is architecture specific
- Base type's ordering can specify both byte and bit ordering
- Works on any location kind
- Locations with bit offsets allowed in composite location parts just like any other location

Source Program	
Variable	
Hardware	$\overline{\mathbf{V}}$
SGPR 3 20 bits	

#### Other benefits of generalizing locations on the stack

- DWARF 5 only supports memory locations on the stack => uses global memory address:
  - DW\_AT\_data\_member\_location => evaluates expression with type instance object address as initial stack value
  - DW\_OP\_push\_object\_address => pushes location of context's program object defined by the attribute
  - DW\_OP\_call\* operations => values can be passed in/out to called DWARF procedure on stack
- Generalization allows any location kind
  - Necessary to support optimized code on GPUs => compiler allocates objects in registers, different address spaces, and composites of them
  - Allows bit fields and implicit locations to be supported => can occur through optimization on any target
- GPU compiler uses DWARF procedures to factorize location expressions => SIMT divergent control flow information
  - Reduces DWARF size
  - More convenient to generate

## **Call Frame Information (CFI)**

- DWARF defines call frame information (CFI) => used to virtually unwind call stack
- Extended CFI rules to support:
  - All location kinds
  - Address spaces
- GPU only saves active lanes of VGPR callee saved registers

DW\_OP\_LLVM\_select\_bit\_piece => used by unwind expressions to inspect the bits in EXEC register DW OP LLVM call frame entry reg => used to get EXEC register value on entry to function

### **Multiple Places**

- DWARF 5 supports loclists => can specify a location is in multiple places at same time
- DW\_OP\_call\* and DW\_OP\_implicit\_pointer => can specify DIE that has a loclist
- Location extended to allow one or more single locations
- Location operations extended to act on multiple places
  - DW\_OP\_LLVM\_offset => adjusts offset of all the single locations
- DWARF 5 defines operation expressions and loclist expressions separately
  - Works in DWARF 5 as locations can only be the last step of an expression
  - Extension generalizations made unification fall out naturally => unification necessary as locations now allowed at any step of an expression

### **Extension Operation Summary**

#### **Core Extensions**

#### • Expression operations:

- DW\_OP\_LLVM\_form\_aspace\_address
- DW\_OP\_LLVM\_push\_lane
- DW\_OP\_LLVM\_offset
- DW\_OP\_LLVM\_offset\_uconst
- DW\_OP\_LLVM\_bit\_offset
- DW\_OP\_LLVM\_call\_frame\_entry\_reg
- DW\_OP\_LLVM\_undefined
- DW\_OP\_LLVM\_aspace\_bregx
- DW\_OP\_LLVM\_aspace\_implicit\_pointer
- DW\_OP\_LLVM\_piece\_end
- DW\_OP\_LLVM\_extend
- DW\_OP\_LLVM\_select\_bit\_piece

#### • CFI operations:

- DW\_CFA\_LLVM\_def\_aspace\_cfa
- DW\_CFA\_LLVM\_def\_aspace\_cfa\_sf

#### • DIE Attributes:

• DW\_AT\_LLVM\_vector\_size

#### **Divergent Lane Support Extensions**

#### • DIE Attributes:

- DW\_AT\_LLVM\_active\_lane
- DW\_AT\_LLVM\_lanes
- DW\_AT\_LLVM\_lane\_pc

## **Current Progress**

- Ongoing development:
  - AMD ROCm ROCgdb debugger <u>https://github.com/ROCm-Developer-Tools/ROCgdb</u>
- In development:
  - AMD ROCm LLVM compiler
  - Perforce TotalView debugger
  - Mentor Graphics GCC compiler
- Further Information:
  - DWARF Extensions For Heterogeneous Debugging <u>https://llvm.org/docs/AMDGPUDwarfExtensionsForHeterogeneousDebugging.html</u>
  - User Guide for AMDGPU Backend: DWARF Debug Information <u>https://llvm.org/docs/AMDGPUUsage.html#dwarf-debug-information</u>

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  - HPE: Andrew Gontarek, Deepak Eachempati, John Vogt, Jeff Sandoval
  - Intel: Markus Metzer
  - Lawrence Livermore National Laboratory: Dong Ahn
  - Mentor Graphics: Andrew Stubbs
  - Oak Ridge National Laboratory
  - Perforce: John DelSignore, Steve Lawrence

## Summary

- DWARF expressions are generalized to allow locations on the stack
- New operators that are composable, consistent, and backward compatible
- Provides support needed by GPUs and other heterogeneous devices
- Improves debugging of optimized code for both CPUs and GPUs

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